

Location Problem for Compressor Stations in Pipeline Networks

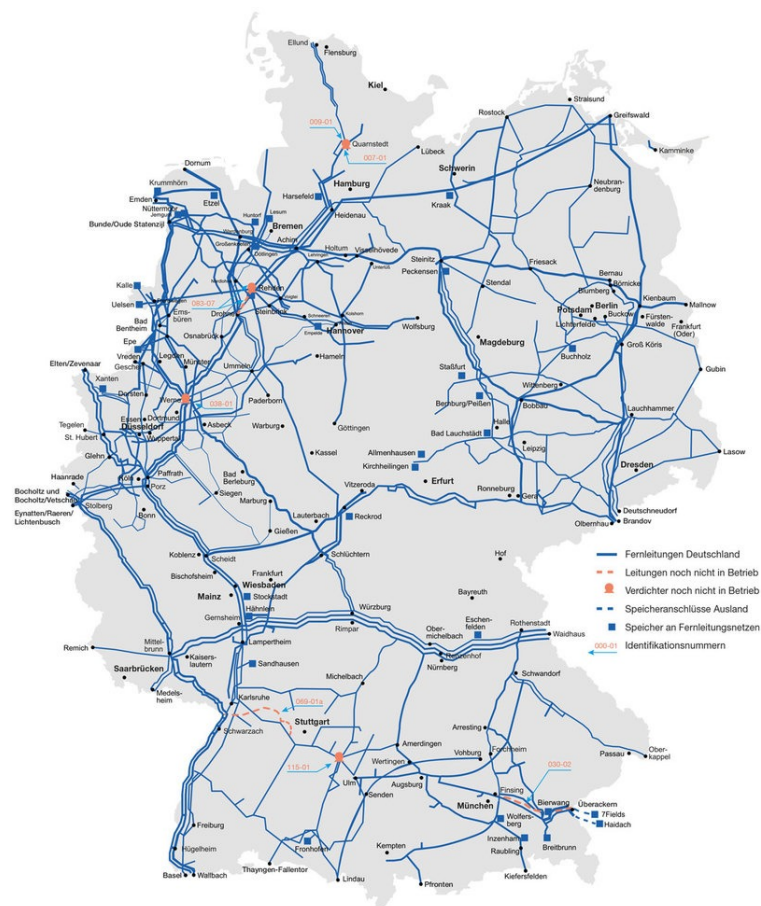
Michael Schuster

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Friedrich-Alexander Universität Erlangen-Nürnberg (FAU),
Chair for Dynamics, Control, Machine Learning and Numerics (FAU DCN-AvH)

Motivation

Natural Gas Transport



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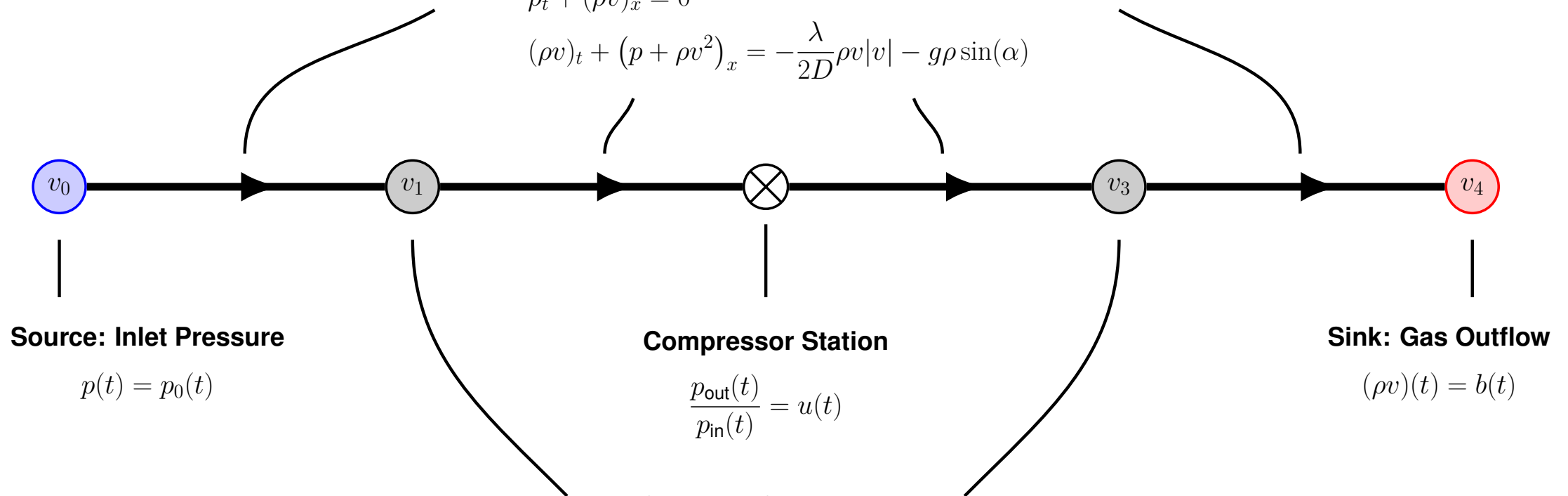
Gas Flow in Pipeline Networks

p	gas pressure	v	gas velocity	g	gravitational constant
ρ	gas density	λ/D	pipe friction	α	pipe slope

Isothermal Euler Equations

$$\rho_t + (\rho v)_x = 0$$

$$(\rho v)_t + (p + \rho v^2)_x = -\frac{\lambda}{2D} \rho v |v| - g \rho \sin(\alpha)$$



Source: Inlet Pressure

$$p(t) = p_0(t)$$

Compressor Station

$$\frac{p_{\text{out}}(t)}{p_{\text{in}}(t)} = u(t)$$

Sink: Gas Outflow

$$(\rho v)(t) = b(t)$$

Coupling Conditions

Conservation of Mass: $\sum (\rho v)_{\text{in}}(t) = \sum (\rho v)_{\text{out}}(t),$ Continuity in Pressure: $p_{\text{in}} = p_{\text{out}}$

Model Hierarchy

(ISO1) - quasilinear model

$$\rho_t + q_x = 0$$
$$q_t + \left(p + \frac{q^2}{\rho} \right)_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

(ISO3) - friction dominated model

$$\rho_t + q_x = 0$$
$$p_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

(ISO2) - semilinear model

$$\rho_t + q_x = 0$$
$$q_t + p_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

(ISO4) - stationary model

$$q_x = 0$$
$$p_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

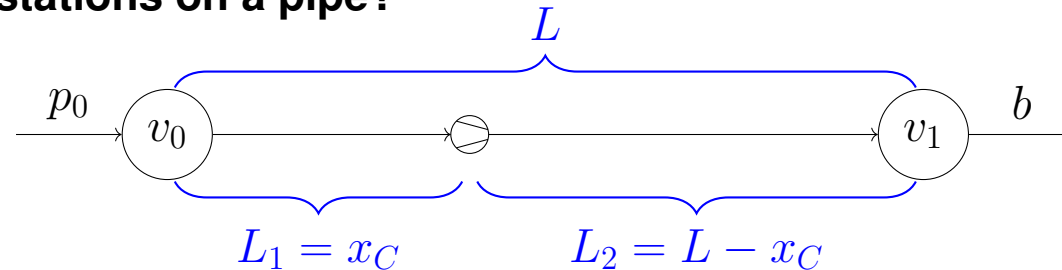
P. Domschke, B. Hiller, J. Lang, C. Tischendorf (2017): *Modellierung von Gasnetzwerken: Eine Übersicht*. Technical Report 2717, Technische Universität Darmstadt, <http://www3.mathematik.tu-darmstadt.de/fb/mathe/preprints.html>

Mathematical Modelling

- The stationary gas flow for ideal gases on a single pipe is given by

$$q(x) \equiv b \quad (\text{const.}), \quad p(x) = p_0^2 - \phi b |b| x \quad \text{with} \quad \phi = \frac{\lambda}{D} R_S T, \quad x \in [0, L]$$

- Where to place compressor stations on a pipe?



- The stationary gas flow with compressor station for ideal gas is given by

$$p_1^2(x) = p_0^2 - \phi b |b| x \quad x \in [0, L_1]$$

$$p_2^2(x) = u p_0^2 - \phi b |b| (u L_1 + x) \quad x \in [0, L_2]$$

- Consider pressure bounds on the pipe

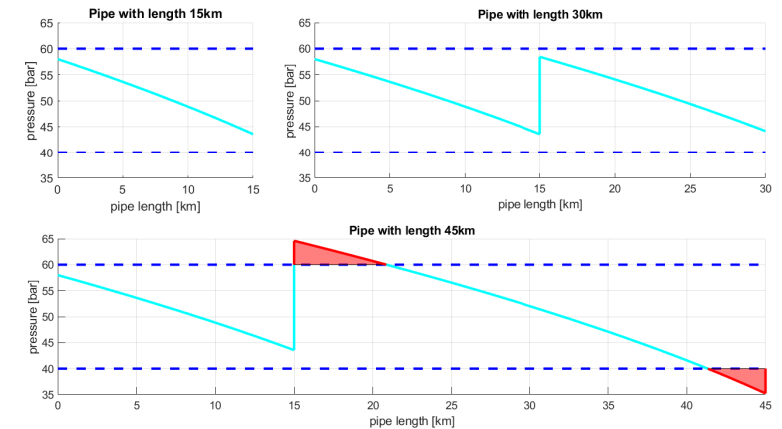
$$\begin{aligned} p_1(x) &\in [p_{\min}, p_{\max}] \\ p_2(x) &\in [p_{\min}, p_{\max}] \end{aligned} \iff \begin{aligned} p_1(0) &\leq p_{\max}, & p_1(L_1) &\geq p_{\min}, \\ p_2(0) &\leq p_{\max}, & p_2(L_2) &\geq p_{\min} \end{aligned}$$

Optimal Compressor Location

Deterministic Optimization

Consider the following optimization problem:

$$\text{(OPT 1)} \quad \left\{ \begin{array}{l} \min_{u, x_C} \quad u^2, \\ \text{s.t.} \quad p_1(L_1) \geq p_{\min}, \quad p_2(0) \leq p_{\max}, \quad p_2(L_2) \geq p_{\min}, \\ \quad \quad u \geq 1, \\ \quad \quad x_C \in [0, L]. \end{array} \right.$$

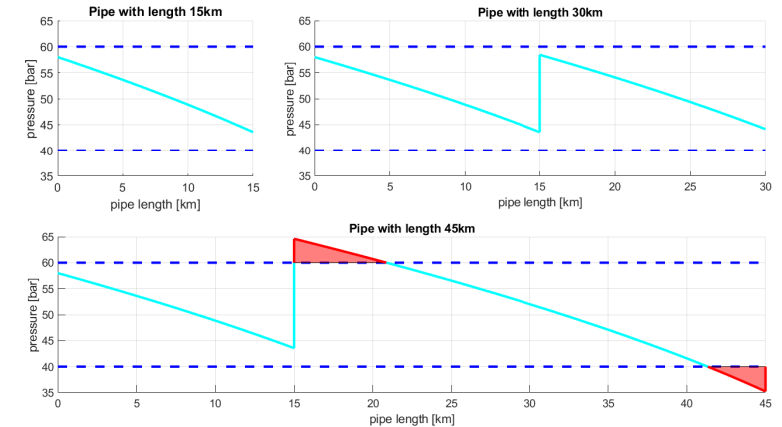


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Lemma

Let $p_0 \in [p_{\min}, p_{\max}]$ and $b > 0$ be given.

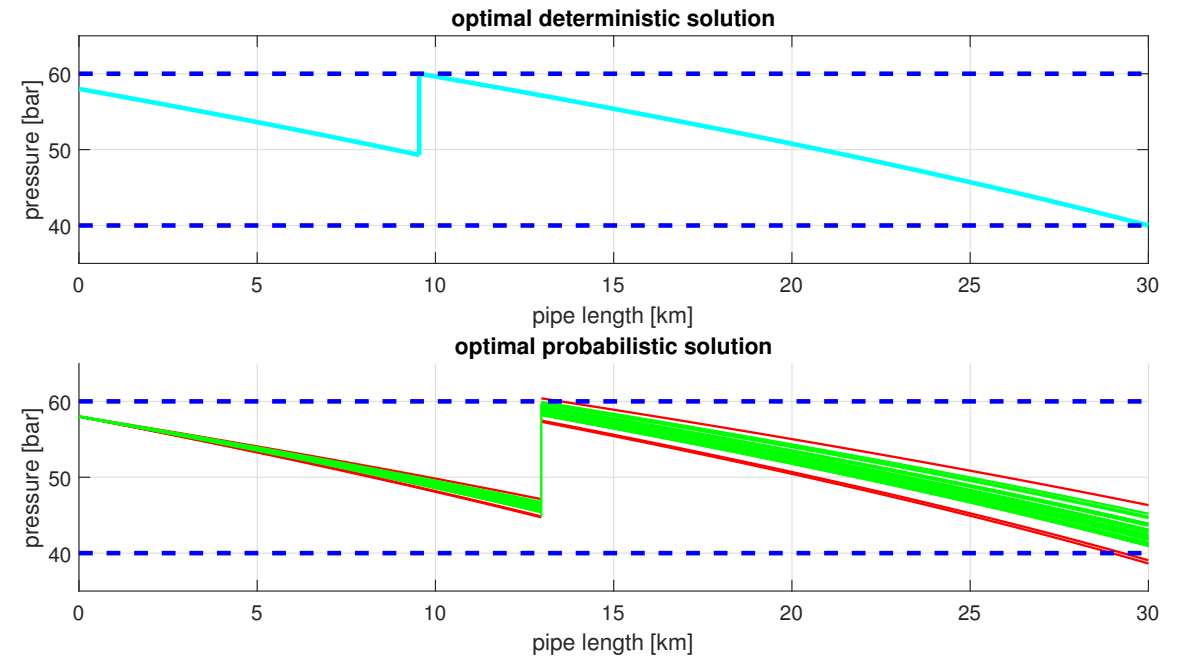
- (i) For $L \leq \frac{p_0^2 - p_{\min}^2}{\phi b |b|}$ every point (u, x_C) with $u = 1$ and $x_C \in [0, L]$ is a solution of the optimization problem (OPT 1).
- (ii) For $\frac{p_0^2 - p_{\min}^2}{\phi b |b|} < L \leq \frac{p_0^2 + p_{\max}^2 - 2 p_{\min}^2}{\phi b |b|}$ the optimization problem (OPT 1) has a unique solution (u^*, x_C^*) with $u^* > 1$ and $x_C \in [0, L]$.
- (iii) For $L > \frac{p_0^2 + p_{\max}^2 - 2 p_{\min}^2}{\phi b |b|}$ the optimization problem (OPT 1) does not have a solution.

Gas outflow b is random in the sense that

$$b = \xi(\omega), \quad \xi \sim \mathcal{N}(\mu, \sigma)$$

Consider the following optimization problem:

$$(OPT\ 2) \quad \left\{ \begin{array}{l} \min_{u, x_C} u^2, \\ \text{s.t.} \quad \mathbb{P} \left(\begin{array}{l} p_1(L_1) \geq p_{\min} \\ p_2(0) \leq p_{\max} \\ p_2(L_2) \geq p_{\min} \\ p_2(L_2) \leq p_{\max} \end{array} \right) \geq \alpha, \\ u \geq 1, \\ x_C \in [0, L]. \end{array} \right.$$



Lemma

If (u^*, x_C^*) with $u^* > 1$ is a solution of (OPT 2), then the probabilistic constraint is active.

Theorem

Let $p_0 \in [p_{\min}, p_{\max}]$ be given.

- (i) If there exists a pair (u, x_C) with $u = 1$ and $x_C \in [0, L]$, that satisfies the constraints of (OPT 2), then every pair (u, x_C) with $u = 1$ and $x_C \in [0, L]$ is a solution of (OPT 2).
- (ii) If there exist a pair (u, x_C) , that satisfies the constraints of (OPT 2) and if (u, x_C) with $u = 1$ is infeasible for at least one $x_C \in [0, L]$, then there exists at least one solution (u^*, x_C^*) of (OPT 2) with $u^* > 1$ and $x_C^* \in [0, L]$.

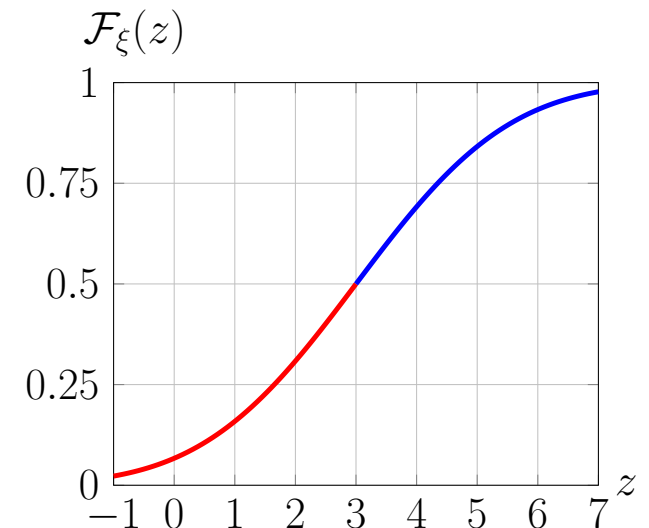
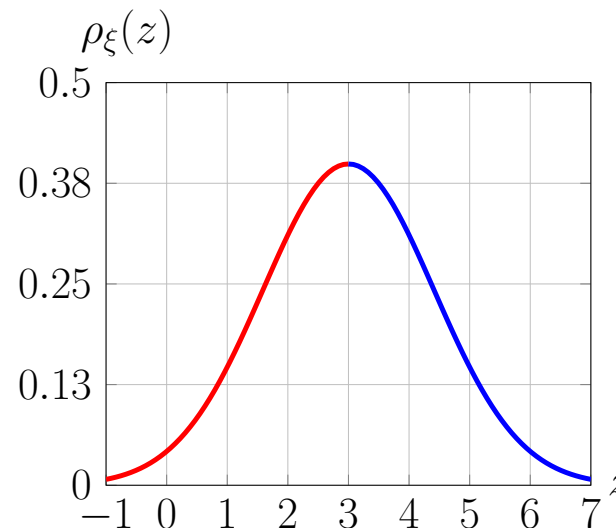
Theorem

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- (ii) If there exist a pair (u, x_C) , that satisfies the constraints of (OPT 2) and if (u, x_C) with $u = 1$ is infeasible for at least one $x_C \in [0, L]$, then there exists at least one solution (u^*, x_C^*) of (OPT 2) with $u^* > 1$ and $x_C^* \in [0, L]$.

Theorem

Let $\alpha > \frac{1}{2}$ be given. For a Gaussian distribution, *Statement (ii) in the last Theorem* guarantees the existence of a unique solution (u^*, x_C^*) of (OPT 2) with $u^* > 1$ and $x_C^* \in [0, L]$.



Optimal Compressor Location

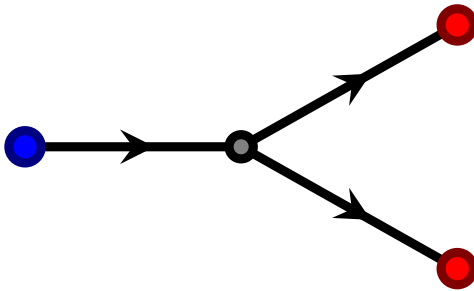
Mathematical Modelling on Networks

- Consider a connected, directed graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$ with vertex set $\mathcal{V}$ and set of edges $\mathcal{E}$
- Binary variables δ_i states if a compressor location is located on edge e_i
- The stationary gas flow for ideal gas on pipe e_i is given by

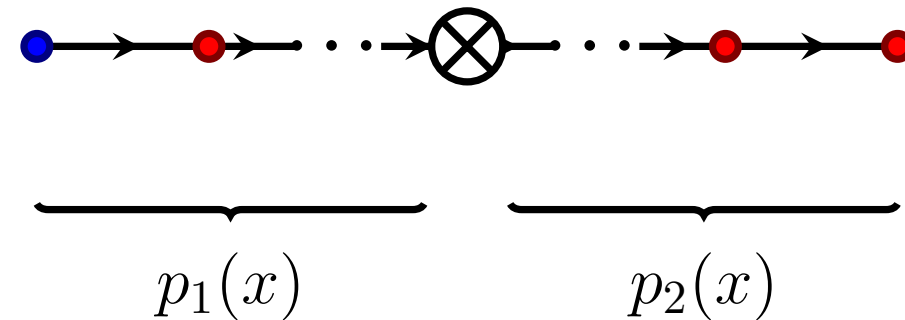
$$p_{i,1}^2(x) = p_{i,1}^2(0) - \phi q_i |q_i| x \quad x \in [0, \delta_i x_{C,i}]$$

$$p_{i,2}^2(x) = (1 - \delta_i + \delta_i u_i) p_{i,1}^2(L_{i,1}) - \phi q_i |q_i| x \quad x \in [0, L - \delta_i x_{C,i}]$$

- Uniqueness in general mainly depends on the graph topology



(a) Scheme of a symmetric graph with one source and two sinks



(b) Scheme of a linear graph with one source and n sinks

Optimal Compressor Location

Deterministic Optimization on Networks

Consider the deterministic optimization problems

$$\text{(OPT 3)} \quad \left\{ \begin{array}{l} \min_{u, x_C, \delta} \quad \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t.} \quad \text{for all } i = 1, \dots, m, \text{ we have} \\ p_{i,1}(L_{i,1}) \geq p_{\min}, \\ p_{i,2}(0) \leq p_{\max}, \\ p_{i,2}(L_{i,2}) \geq p_{\min}, \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C, \end{array} \right.$$

$$\text{(OPT 4)} \quad \left\{ \begin{array}{l} \min_{u, x_C, \delta, n_C} \quad \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t.} \quad \text{for all } i = 1, \dots, m, \text{ we have} \\ p_{i,1}(L_{i,1}) \geq p_{\min}, \\ p_{i,2}(0) \leq p_{\max}, \\ p_{i,2}(L_{i,2}) \geq p_{\min}, \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C. \end{array} \right.$$

Optimal Compressor Location

Deterministic Optimization on Networks

Consider the deterministic optimization problems

$$(OPT\ 3) \left\{ \begin{array}{l} \min_{u, x_C, \delta} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. for all } i = 1, \dots, m, \text{ we have} \\ p_{i,1}(L_{i,1}) \geq p_{\min}, \\ p_{i,2}(0) \leq p_{\max}, \\ p_{i,2}(L_{i,2}) \geq p_{\min}, \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C, \end{array} \right.$$

$$(OPT\ 4) \left\{ \begin{array}{l} \min_{u, x_C, \delta, n_C} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. for all } i = 1, \dots, m, \text{ we have} \\ p_{i,1}(L_{i,1}) \geq p_{\min}, \\ p_{i,2}(0) \leq p_{\max}, \\ p_{i,2}(L_{i,2}) \geq p_{\min}, \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C. \end{array} \right.$$

Theorem

For all $v_i \in \mathcal{V}_{\text{in}}$ let $p_{i,0} \in [p_{\min}, p_{\max}]$ be given and for all $v_i \in \mathcal{V}_{\text{out}}$ let $b_i \geq 0$ be given. Further let a number $n_C \in \{0, \dots, m\}$ be given.

- (i) If a triple $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$ with $u_j = 1$ for all $j = 1, \dots, m$ satisfies the constraints of (OPT 3), every triple $(\mathbf{1}_m, x_C, \delta)$ with $x_{C,j} \in [0, L_j]$, $\delta_j \in \{0, 1\}$ and $\sum_{j=1}^m \delta_j = n_C$ is a solution of (OPT 3).
- (ii) If there exists a triple $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$ that satisfies the constraints in (OPT 3), and if $(\mathbf{1}_m, x_C, \delta)$ is infeasible for at least one pair (x_C, δ) with $x_{C,j} \in [0, L_j]$ and $\delta_j \in \{0, 1\}$, the optimization problem (OPT 3) has at least one solution.

Consider the probabilistic optimization problems

$$\text{(OPT 6)} \quad \left\{ \begin{array}{l} \min_{u, x_C, \delta} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. } \mathbb{P} \left(\begin{array}{l} p_{k,1}(L_{k,1}) \geq p_{\min} \\ p_{k,2}(0) \leq p_{\max} \\ p_{k,2}(L_{k,2}) \geq p_{\min} \\ p_{k,2}(L_{k,2}) \leq p_{\max} \end{array} \quad \forall k = 1, \dots, m \right) \geq \alpha, \\ \text{and for all } i = 1, \dots, m, \text{ we have} \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C, \end{array} \right.$$

$$\text{(OPT 7)} \quad \left\{ \begin{array}{l} \min_{u, x_C, \delta, n_C} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. } \mathbb{P} \left(\begin{array}{l} p_{k,1}(L_{k,1}) \geq p_{\min} \\ p_{k,2}(0) \leq p_{\max} \\ p_{k,2}(L_{k,2}) \geq p_{\min} \\ p_{k,2}(L_{k,2}) \leq p_{\max} \end{array} \quad \forall k = 1, \dots, m \right) \geq \alpha, \\ \text{and for all } i = 1, \dots, m, \text{ we have} \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C. \end{array} \right.$$

Consider the probabilistic optimization problems

$$\begin{aligned}
 \text{(OPT 6)} \quad & \left\{ \begin{array}{l} \min_{u, x_C, \delta} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. } \mathbb{P} \left(\begin{array}{l} p_{k,1}(L_{k,1}) \geq p_{\min} \\ p_{k,2}(0) \leq p_{\max} \\ p_{k,2}(L_{k,2}) \geq p_{\min} \\ p_{k,2}(L_{k,2}) \leq p_{\max} \end{array} \quad \forall k = 1, \dots, m \right) \geq \alpha, \\ \text{and for all } i = 1, \dots, m, \text{ we have} \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C, \end{array} \right. \\
 \text{(OPT 7)} \quad & \left\{ \begin{array}{l} \min_{u, x_C, \delta, n_C} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. } \mathbb{P} \left(\begin{array}{l} p_{k,1}(L_{k,1}) \geq p_{\min} \\ p_{k,2}(0) \leq p_{\max} \\ p_{k,2}(L_{k,2}) \geq p_{\min} \\ p_{k,2}(L_{k,2}) \leq p_{\max} \end{array} \quad \forall k = 1, \dots, m \right) \geq \alpha, \\ \text{and for all } i = 1, \dots, m, \text{ we have} \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C. \end{array} \right.
 \end{aligned}$$

Theorem

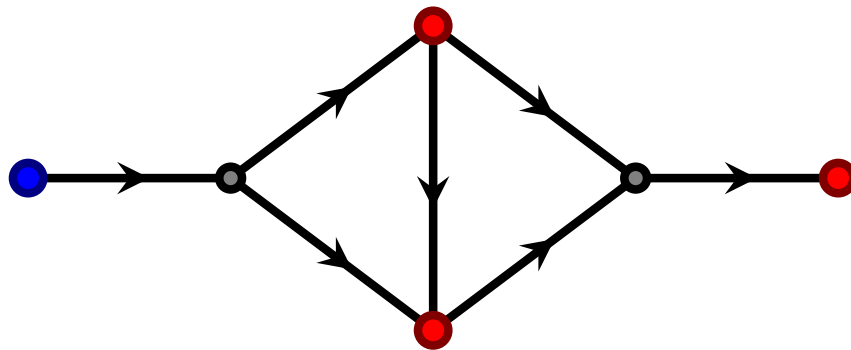
Let $p_{i,0} \in [p_{\min}, p_{\max}]$ be given for every node $v_i \in \mathcal{V}_{\text{in}}$. Further let a number $n_C \in \{0, \dots, m\}$ be given.

- (i) If there exists a triple $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$ with $u = \mathbf{1}_m$, that satisfies the constraints of (OPT 6), then every triple $(\mathbf{1}_m, x_C, \delta)$ with $x_{C,j} \in [0, L_j]$, $\delta_j \in \{0, 1\}$ and $\sum_{j=1}^m \delta_j = n_C$ is a solution of (OPT 6).
- (ii) If there exists a triple $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$, that satisfies the constraints of (OPT 6), and if $(\mathbf{1}_m, x_C, \delta)$ is infeasible for at least one pair (x_C, δ) with $x_{C,j} \in [0, L_j]$ and $\delta \in \{0, 1\}$, then (OPT 6) has at least one solution.

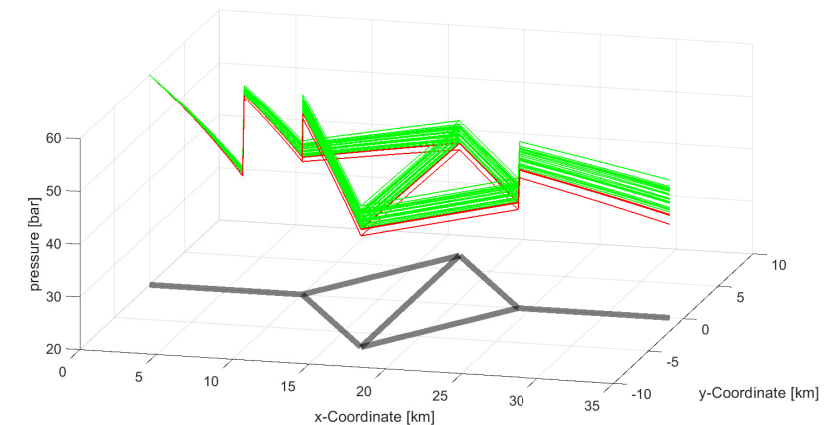
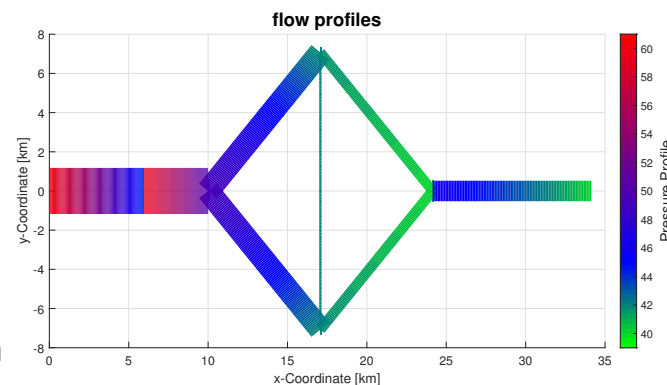
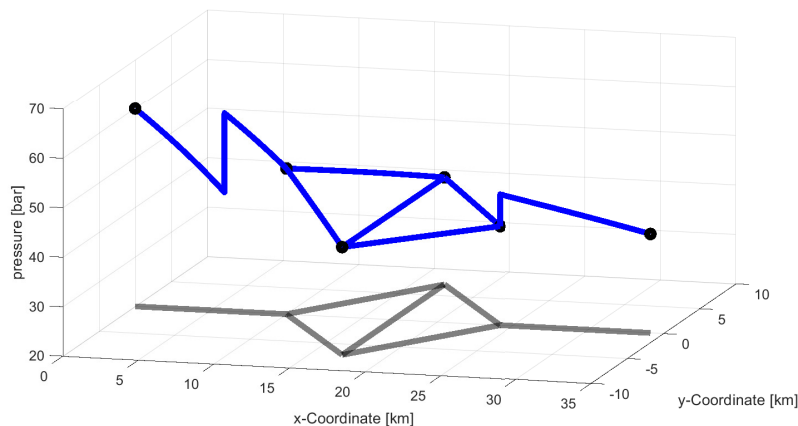
Optimal Compressor Location

A Numerical Example on a Diamond Graph

A scheme of a diamond graph with 1 source (blue) and 3 sinks (red):



Variable	Letter	Value	Unit
inlet pressure	p_0	60	bar
lower pressure bound	$p_{\min}$	40	bar
upper pressure bound	$p_{\max}$	60	bar
gas outflow (=mean value)	$b (= \mu)$	[90, 60, 120]	kg/m ² s
covariance matrix	Σ	diag(2.25, 2.25, 2.25)	
speed of sound in the gas	c	343	m/s
pipe friction coefficient	λ^F	0.1	
pipe diameter	D	0.5	m
specific gas constant	R_S	515	J/kg K
gas temperature	T	293	K
probability level	α	0.8	

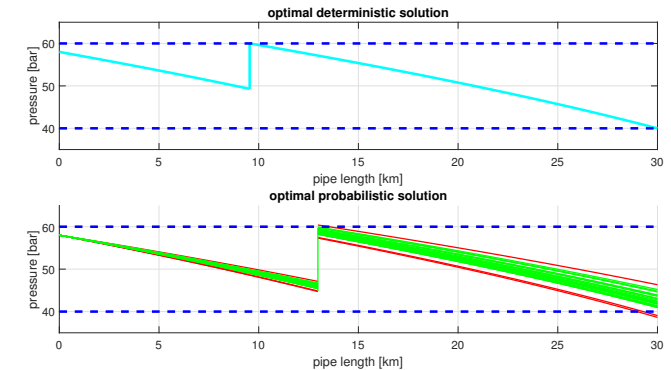


Optimal Compressor Location

A Numerical Example for Transient Gas Flow

- Consider the probabilistic example for the stationary flow on a single pipe
- Solve the probabilistic constraint optimization problem for $\alpha = 90\%$

$$\begin{aligned}
 \text{(OPT 2)} \quad & \left\{ \begin{array}{l} \min_{u, x_C} u^2, \\ \text{s.t. } \mathbb{P} \left(\begin{array}{l} p_1(L_1) \geq p_{\min} \\ p_2(0) \leq p_{\max} \\ p_2(L_2) \geq p_{\min} \\ p_2(L_2) \leq p_{\max} \end{array} \right) \geq \alpha \\ u \geq 1, \\ x_C \in [0, L]. \end{array} \right.
 \end{aligned}$$

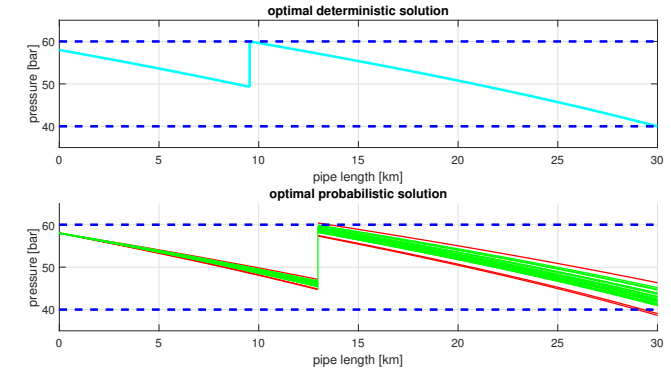


Optimal Compressor Location

A Numerical Example for Transient Gas Flow

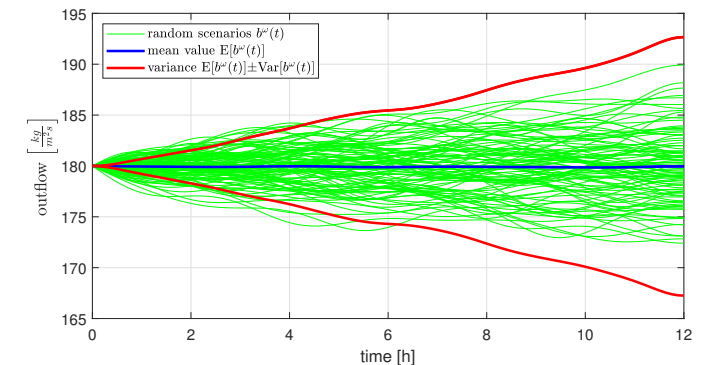
- Consider the probabilistic example for the stationary flow on a single pipe
- Solve the probabilistic constraint optimization problem for $\alpha = 90\%$

$$(\text{OPT } 2) \quad \left\{ \begin{array}{l} \min_{u, x_C} u^2, \\ \text{s.t. } \mathbb{P} \left(\begin{array}{l} p_1(L_1) \geq p_{\min} \\ p_2(0) \leq p_{\max} \\ p_2(L_2) \geq p_{\min} \\ p_2(L_2) \leq p_{\max} \end{array} \right) \geq \alpha \\ u \geq 1, \\ x_C \in [0, L]. \end{array} \right.$$



- Randomize the boundary data in time by a *Wiener process*
- The probabilistic robustness of the steady state control is

$$\mathbb{P} \left(\begin{array}{l} p_1(t, L_1) \geq p_{\min} \\ p_2(t, 0) \leq p_{\max} \\ p_2(t, L_2) \geq p_{\min} \\ p_2(t, L_2) \leq p_{\max} \end{array} \quad \forall t \in [0, T] \right) = 85.74\%$$



Optimal Compressor Location



A Numerical Example for Transient Gas Flow

References

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