

Fault-tolerance by construction

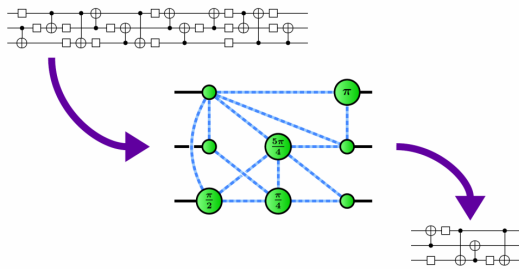
Aleks Kissinger

Fault-tolerant Quantum Technologies, Benasque 2026



ZX-calculus

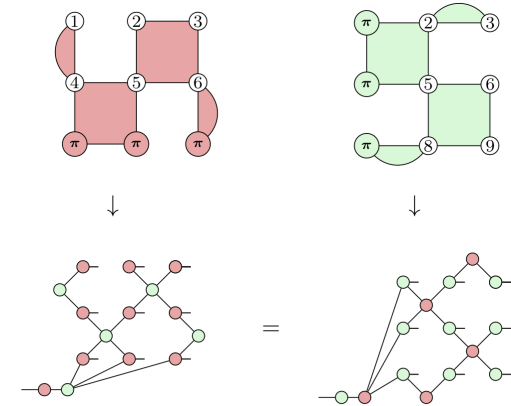
ZX := a handy tool for working with quantum computations using **graph rewriting**



Optimisation

$$\begin{aligned}
 e^{i\pi/4} \begin{array}{c} | \\ \pi/4 \end{array} &= +2e^{i\pi/4} \begin{array}{c} \text{---} \\ \pi/2 \end{array} \\
 -\frac{1+\sqrt{2}}{4} \begin{array}{c} | \\ \pi/4 \end{array} &+ \frac{1-\sqrt{2}}{4} \begin{array}{c} | \\ \pi \end{array} \\
 -2\sqrt{2}i \begin{array}{c} | \\ \pi/2 \end{array} &-2i \begin{array}{c} | \\ \pi \end{array} \\
 +8\sqrt{2}i \begin{array}{c} | \\ \pi/2 \end{array} &+8\sqrt{2}i \begin{array}{c} | \\ \pi \end{array}
 \end{aligned}$$

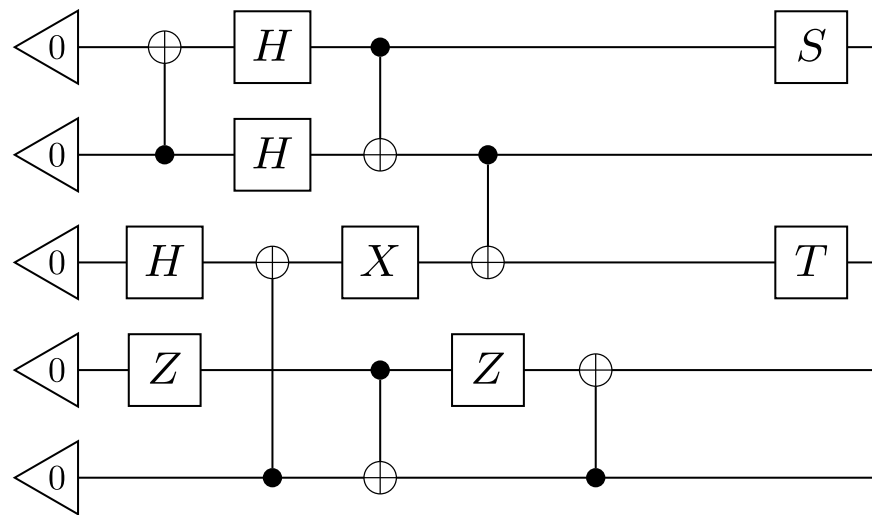
Simulation & Verification

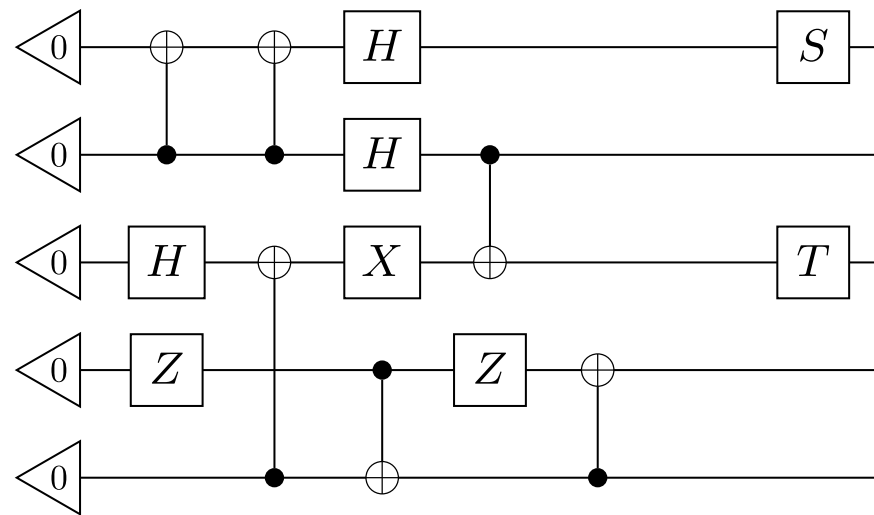


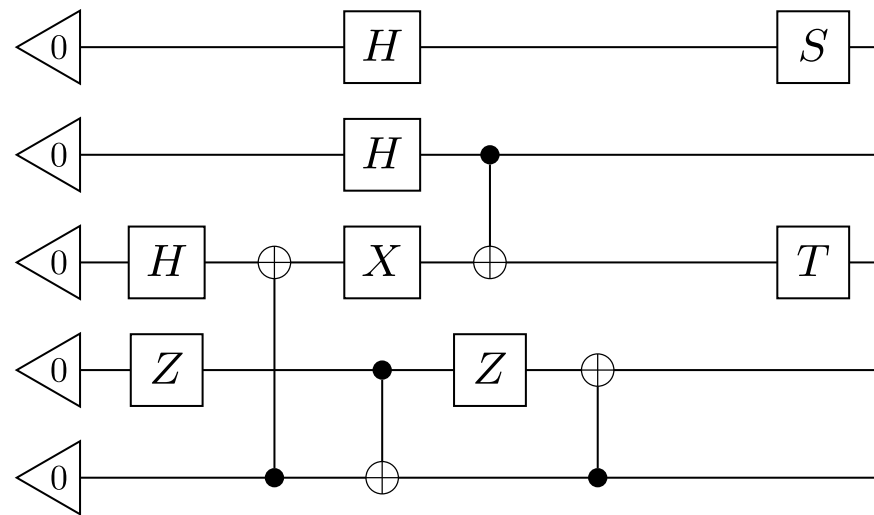
Error Correction

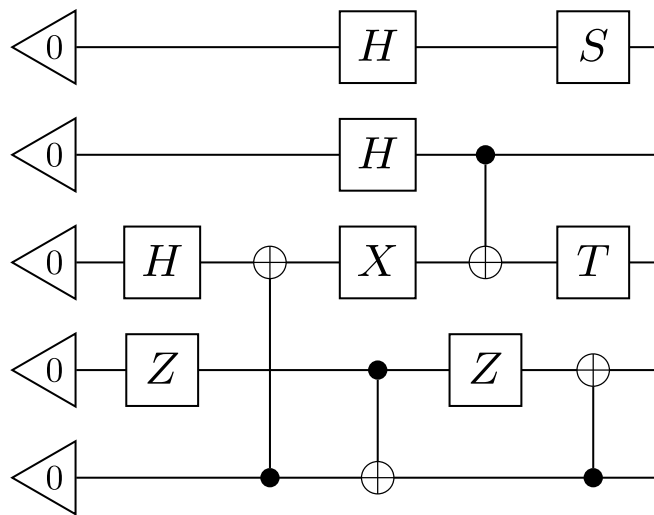
Optimising circuits the old-fashioned way



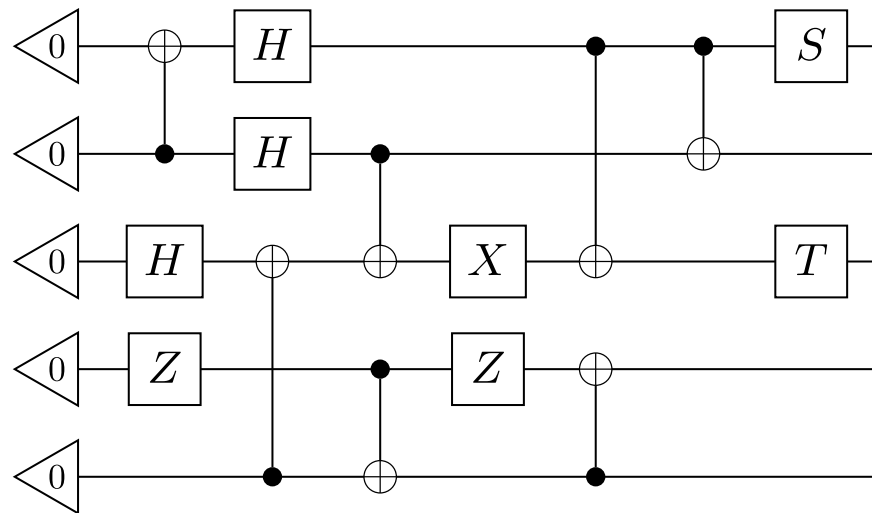




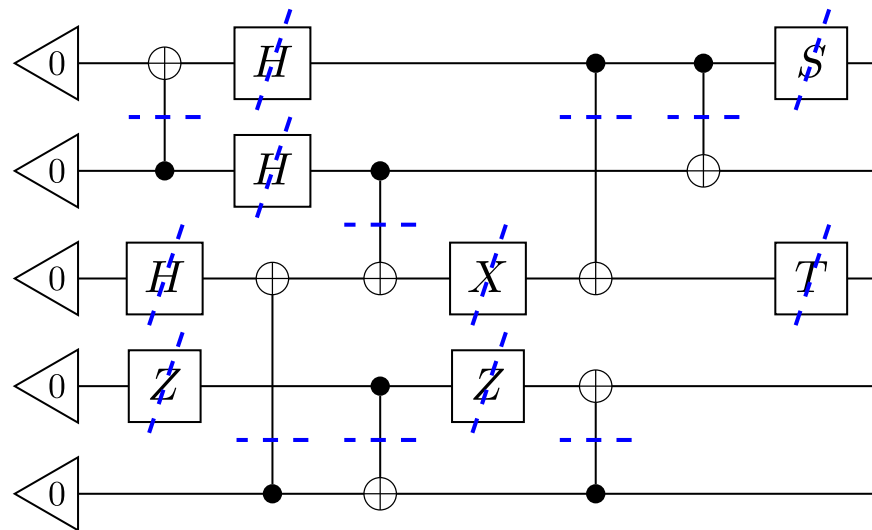




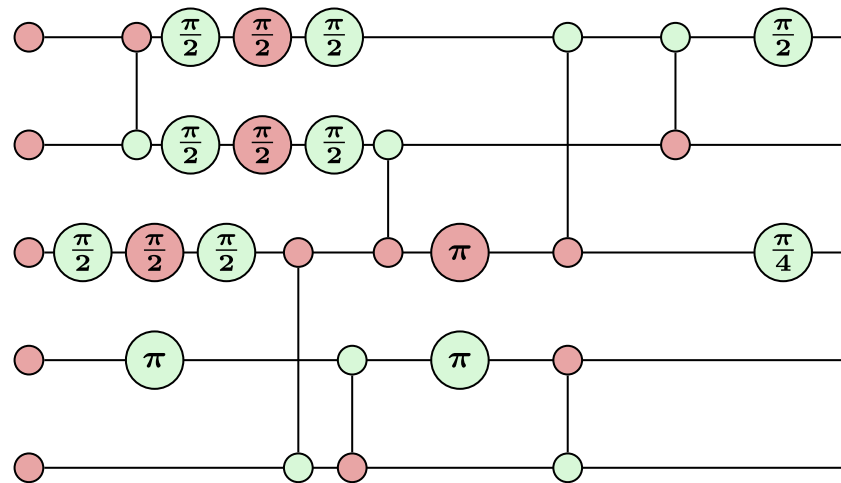
A better idea: decompose a circuit into a ZX diagram



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ZX diagrams

- Gates are represented using more basic building blocks, called **spiders**:

$$\begin{array}{c} \vdots \\ \vdots \end{array} \begin{array}{c} \curvearrowright \\ \alpha \\ \curvearrowleft \end{array} \begin{array}{c} \vdots \\ \vdots \end{array} := |0\dots 0\rangle\langle 0\dots 0| + e^{i\alpha}|1\dots 1\rangle\langle 1\dots 1|$$

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- E.g.

$$CNOT := \begin{array}{c} \text{---} \circ \text{---} \\ | \\ \text{---} \circ \text{---} \end{array} \quad \sqrt{X} := \text{---} \circ \frac{\pi}{2} \text{---} \quad Z_{\alpha} := \text{---} \circ \alpha \text{---}$$

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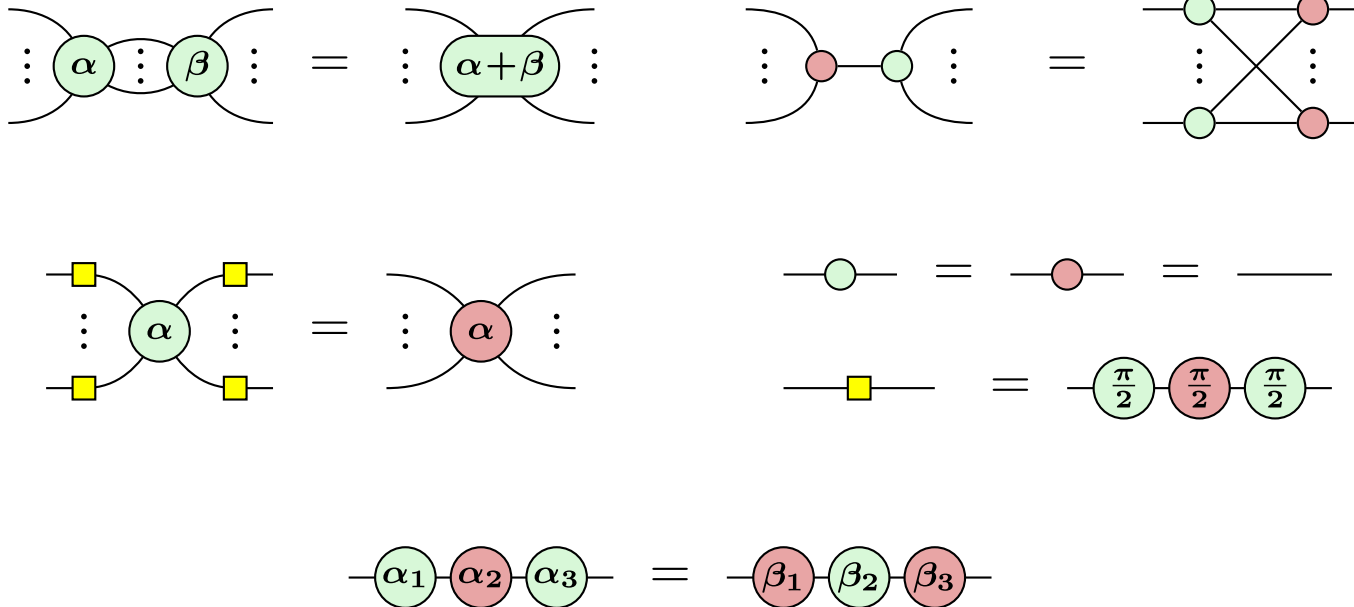
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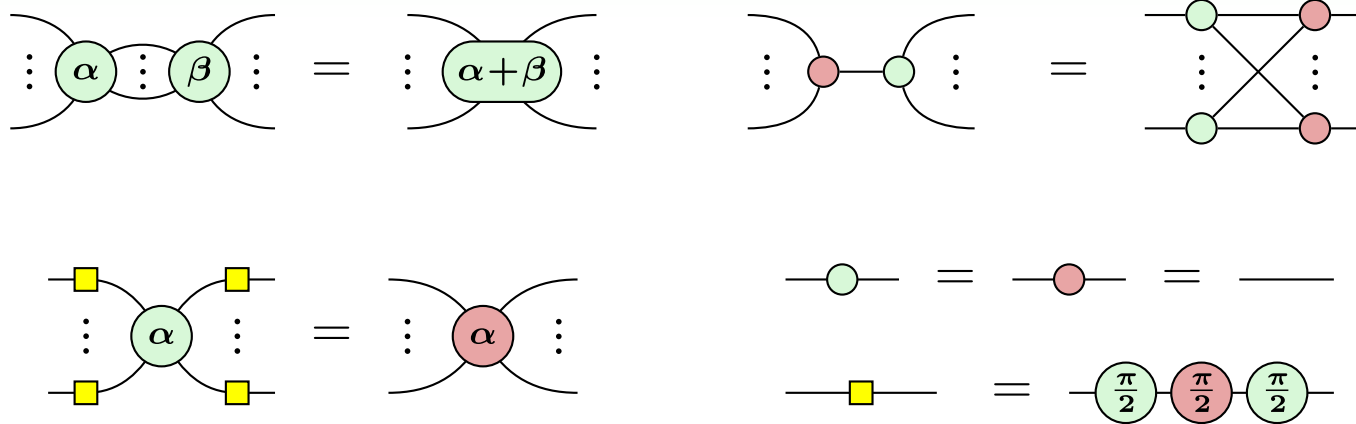
- Wire order doesn't matter $\Rightarrow$ treat ZX-diagrams as **undirected graphs**

...then use the ZX calculus



A complete set of equations for qubit QC

Clifford ZX calculus

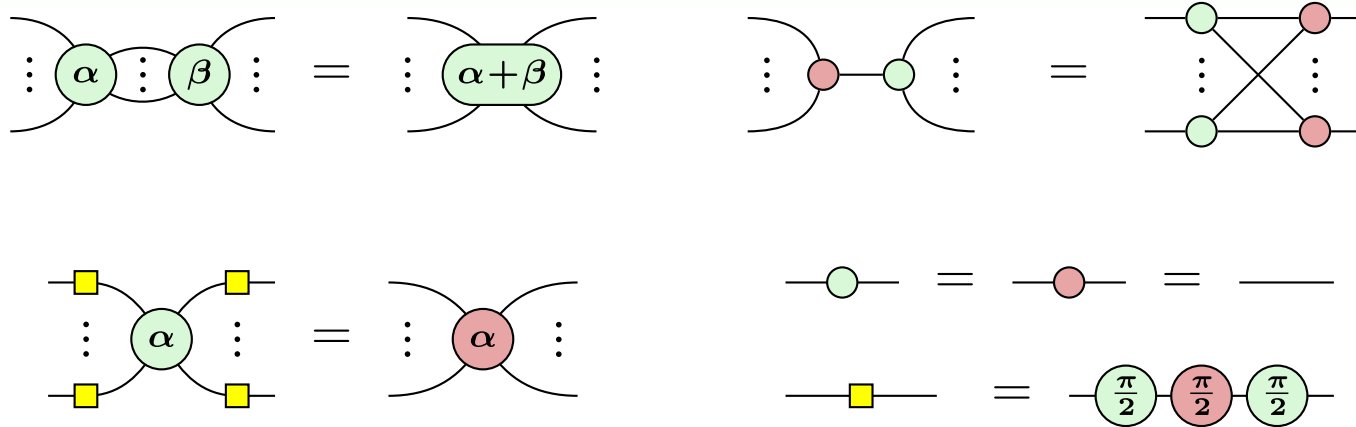


$$\alpha, \beta \in \{0, \frac{\pi}{2}, \pi, -\frac{\pi}{2}\}$$

A complete set of equations for qubit **Clifford** QC

efficient synthesis, equality checking, classical simulation, ...

Clifford ZX calculus

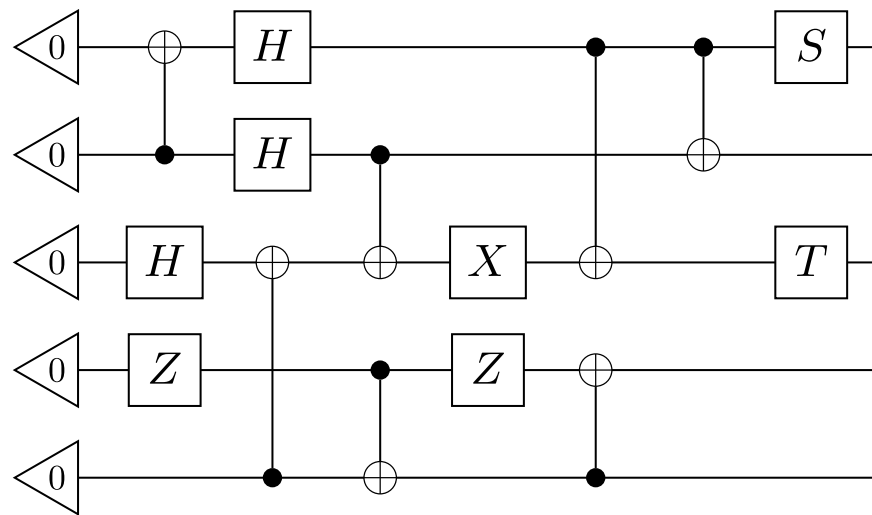


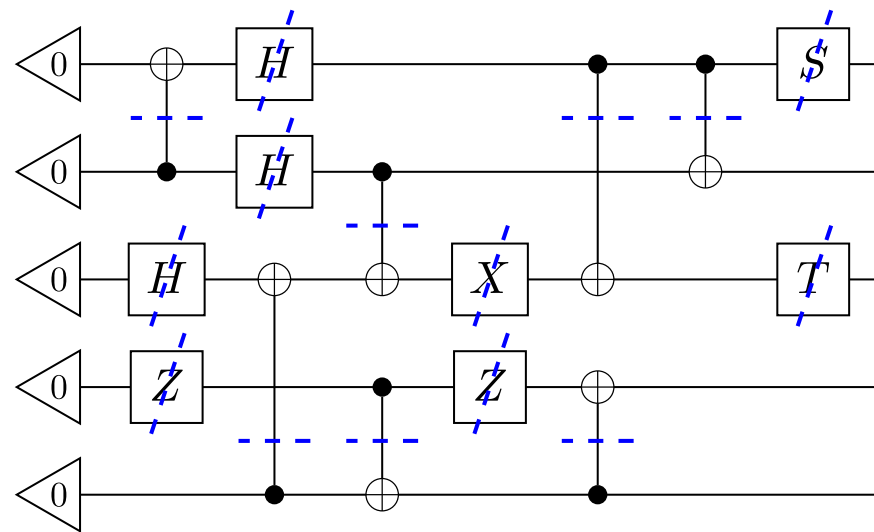
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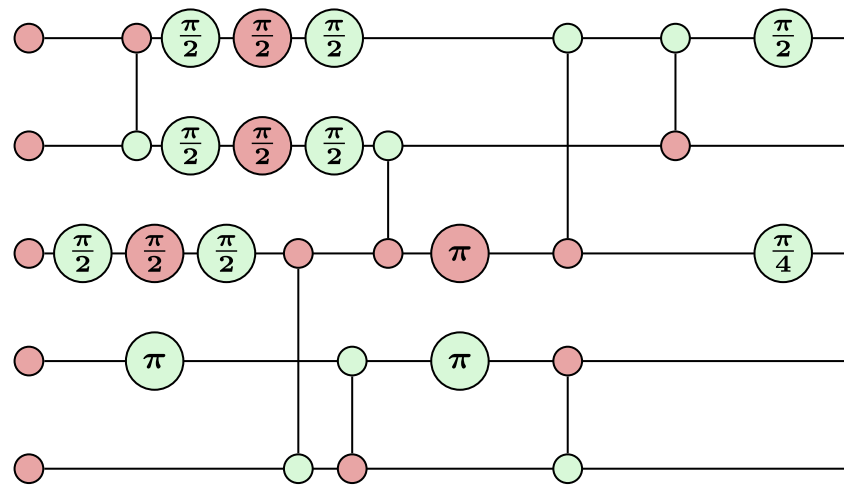
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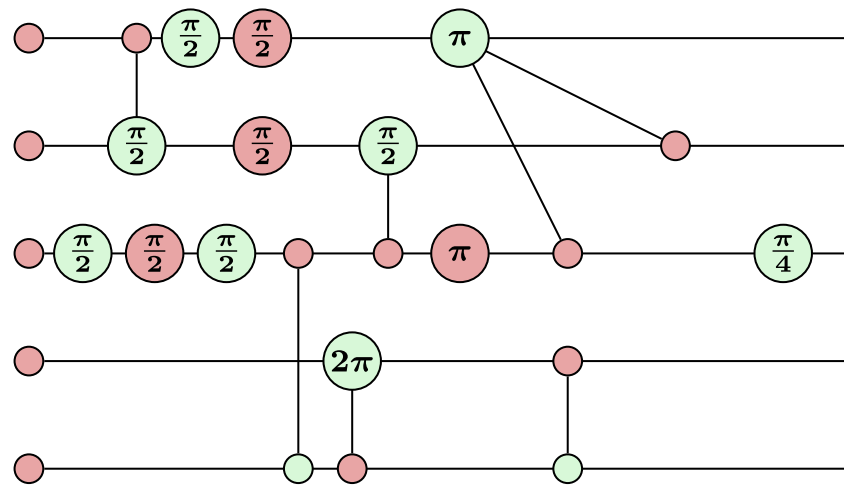
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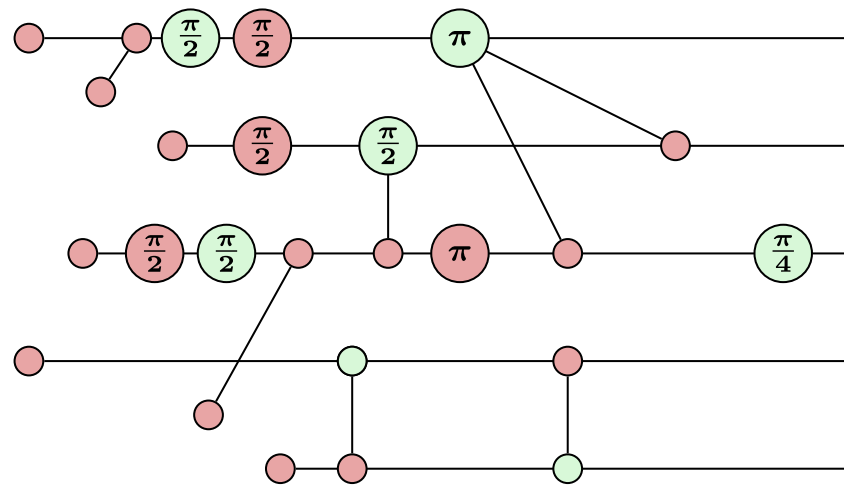
$\approx$ "graphical **stabiliser theory**"

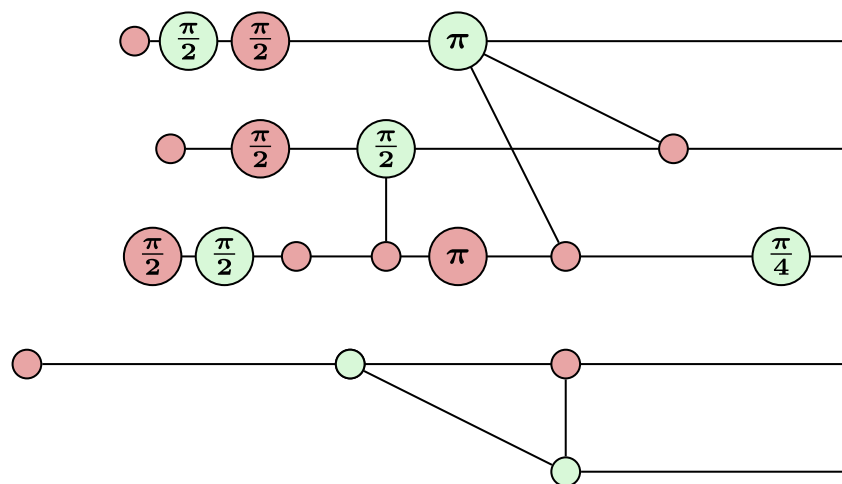


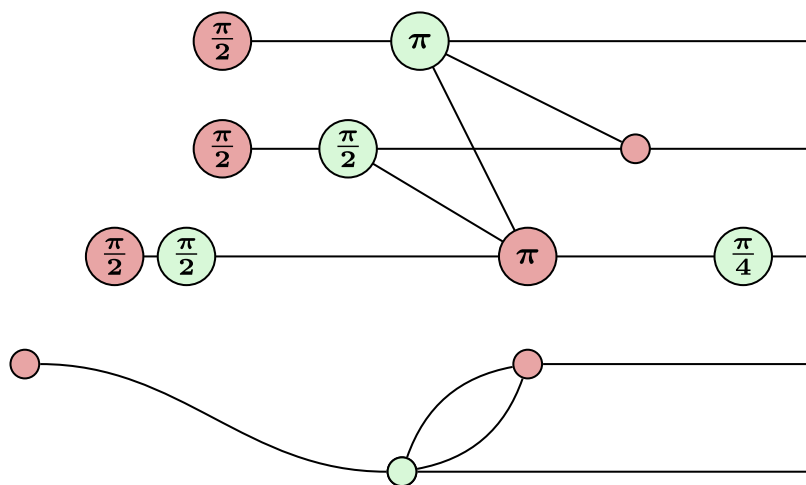


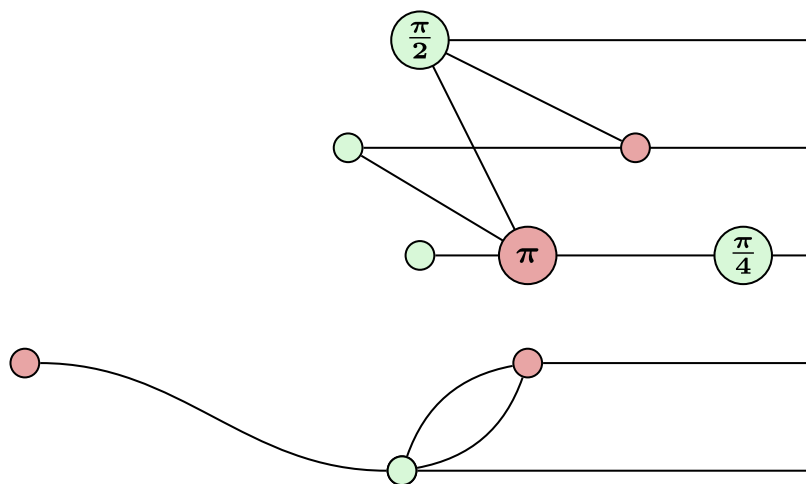


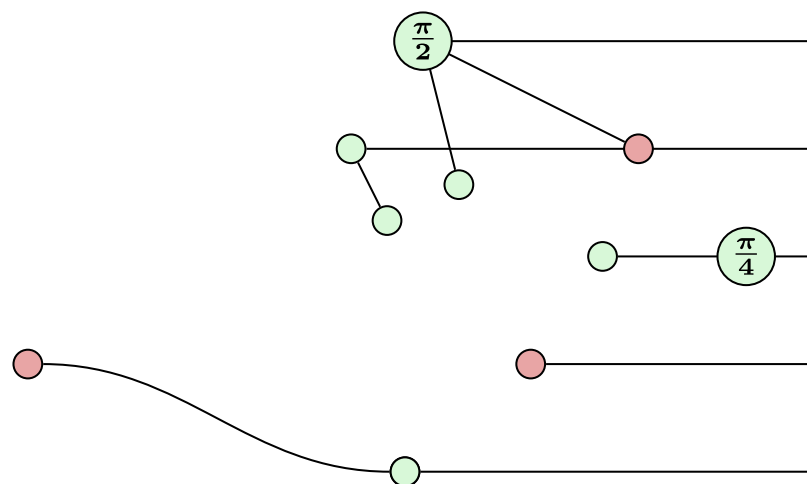


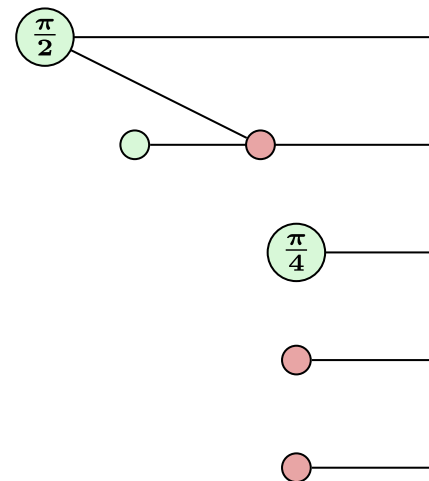




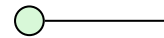




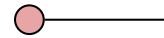
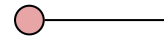




$$\frac{\pi}{2}$$

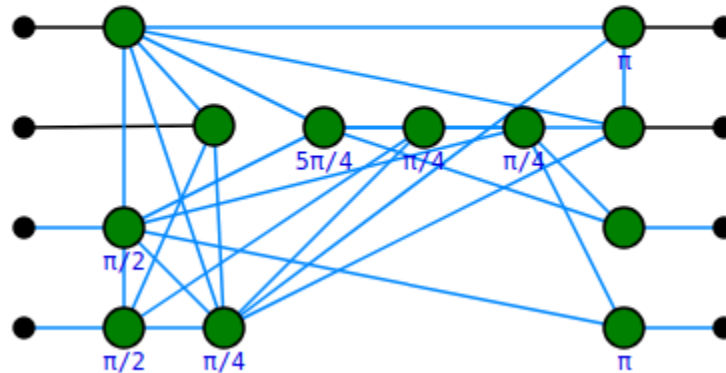


$$\frac{\pi}{4}$$



PyZX

```
In [11]: zx.simplify.clifford_simp(g)
          g.normalise()
          zx.d3.draw(g)
```



- Open source Python library for circuit optimisation, experimentation, and education using ZX-calculus

<https://github.com/zxcalc/pyzx>

QuiZX



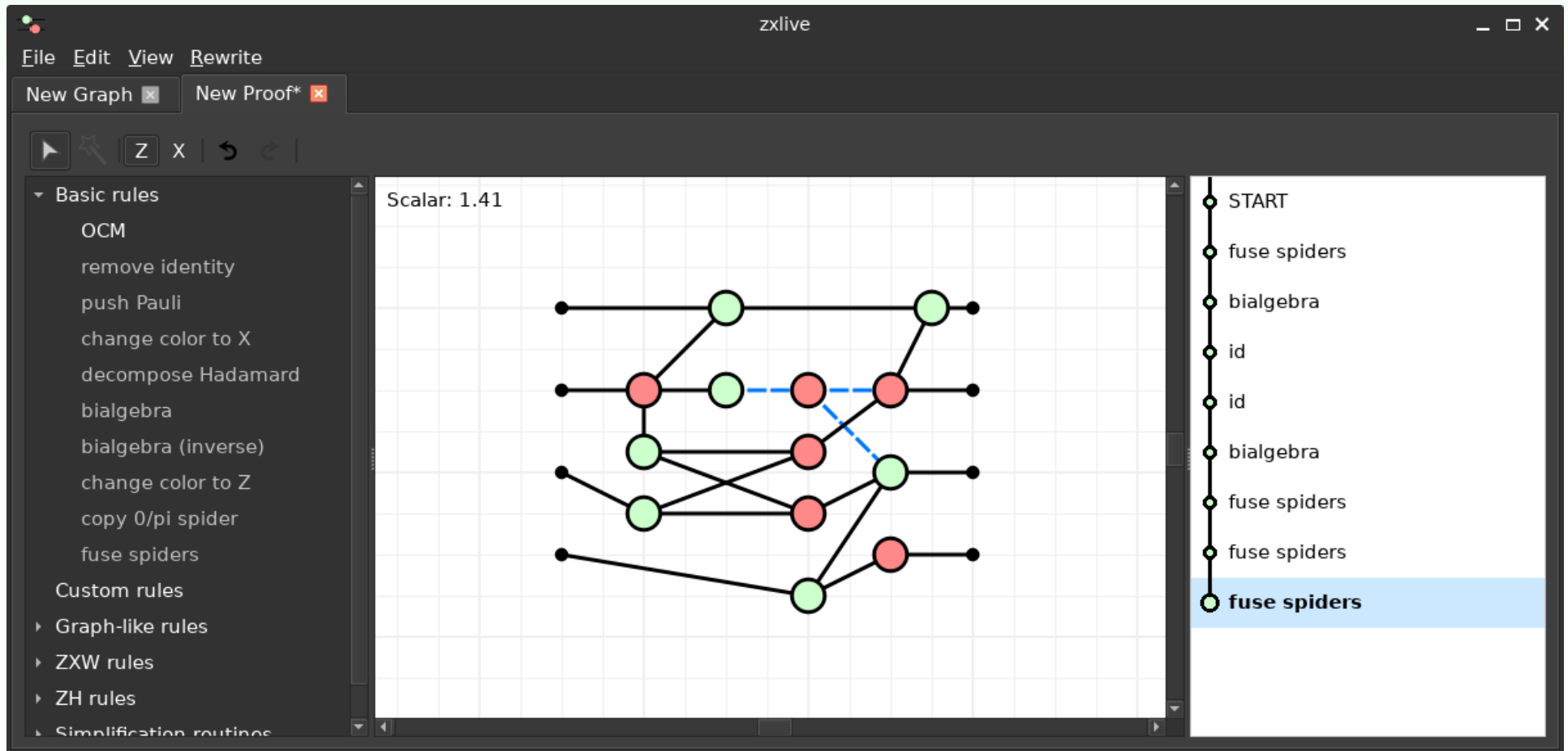
```
use quizx::basic_rules::*;

loop {
    match g.find_edge(|v0,v1,_| check_spider_fusion(&g, v0, v1)) {
        Some((v0,v1,_)) => spider_fusion_unsafe(&mut g, v0, v1),
        None => break,
    };
}
```

- **Large scale** circuit optimisation and classical simulation library for ZX-calculus

<https://github.com/zxcalc/quizx>

ZXLive



- GUI tool based on PyZX

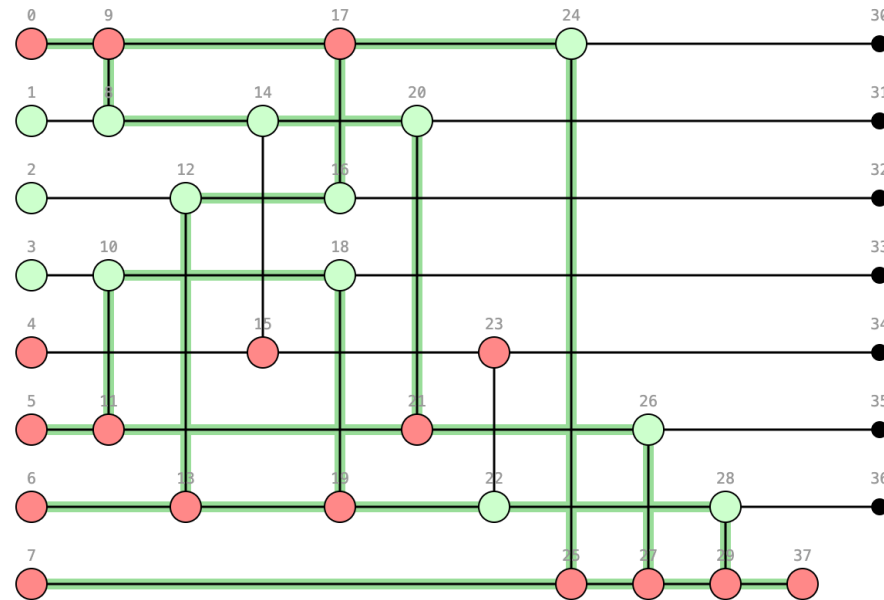
<https://github.com/zxcalc/zxlive>

Paritea

```
from paritea.equivalence import is_fault_equivalence

stabilisers, regions = compute_pauli_webs(d_circuit)
draw(d_circuit, web=regions[0])
is_fault_equivalence(nm_state, nm_circuit, quiet=True)
```

[17] ✓ 0.0s



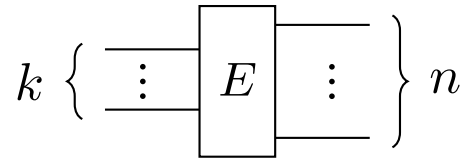
```
... True
```

FT Simulation/Rewriting Utilities: github.com/paritea/paritea

Quantum error correction

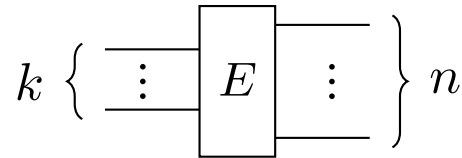
Quantum error correction

...is done by encoding some space of **logical** qubits into a bigger space of **physical** qubits:



Quantum error correction

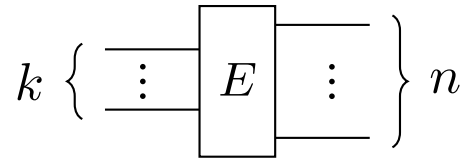
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- E (or just $\text{Im}(E)$) defines a **quantum error correcting code**

Quantum error correction

...is done by encoding some space of **logical** qubits into a bigger space of **physical** qubits:



- E (or just $\text{Im}(E)$) defines a **quantum error correcting code**
- Fault-tolerant quantum computing (FTQC) requires methods for:
 - encoding/decoding logical states and measurements
 - measuring physical qubits to detect/correct errors
 - doing fault-tolerant computations on encoded qubits

QEC in ZX 101...

Quantum Gates

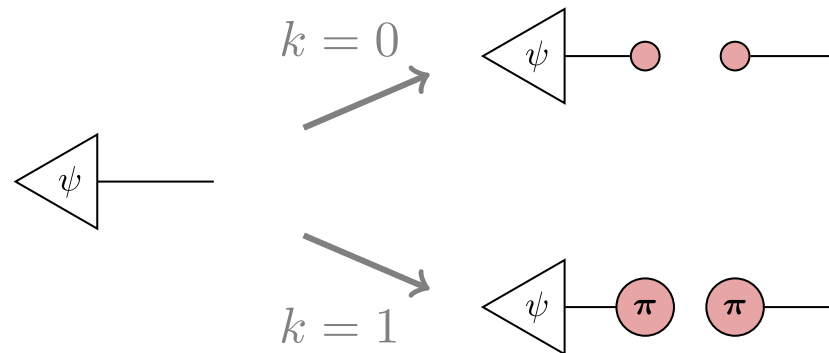
$$CNOT = \begin{array}{c} \text{---} \text{---} \text{---} \\ | \\ \text{---} \text{---} \end{array} \quad CZ = \begin{array}{c} \text{---} \text{---} \text{---} \\ | \\ \text{---} \text{---} \end{array} \quad H = \text{---} \text{---} \text{---} := \text{---} \text{---} \text{---}$$

$$R_Z[\alpha] = \text{---} \text{---} \quad R_X[\alpha] = \text{---} \text{---}$$

Quantum Measurements

$$|k\rangle\langle k| \propto \text{---} \textcircled{k\pi} \textcircled{k\pi} \text{---}$$

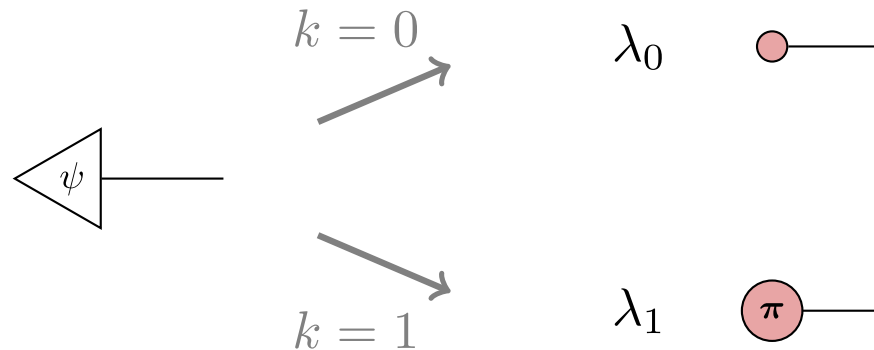
...collapse the quantum state to a fixed one,
depending on the **outcome** $k \in \{0, 1\}$:



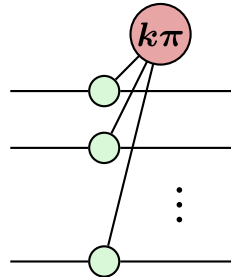
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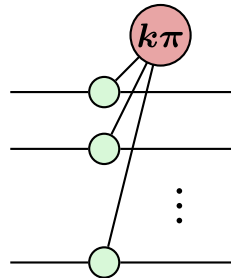
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Multi-qubit measurements



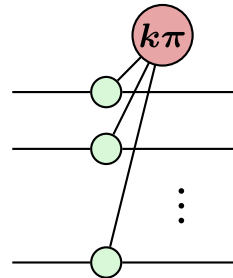
Multi-qubit measurements



$k = 0 \quad \Rightarrow \quad +1$ eigenspace of $Z \otimes \dots \otimes Z$

$k = 1 \quad \Rightarrow \quad -1$ eigenspace of $Z \otimes \dots \otimes Z$

Multi-qubit measurements



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Generalises the single-qubit case, thanks to the "copy" rule:

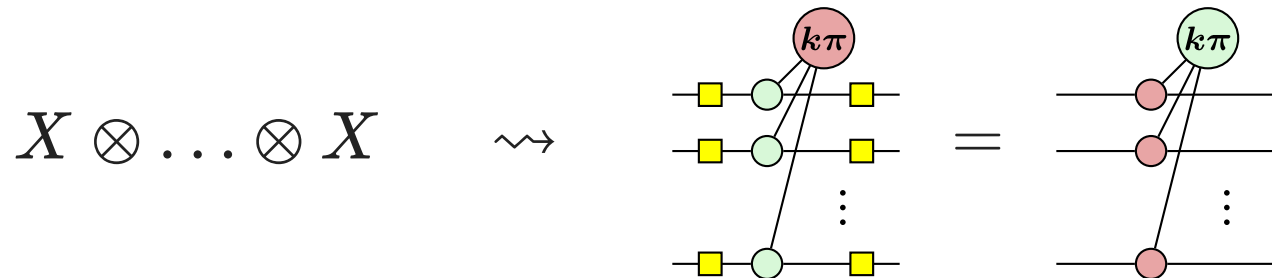


Multi-qubit measurements

Any other multi-qubit measurement can be obtained by conjugating with local Cliffords, e.g.

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Modelling errors

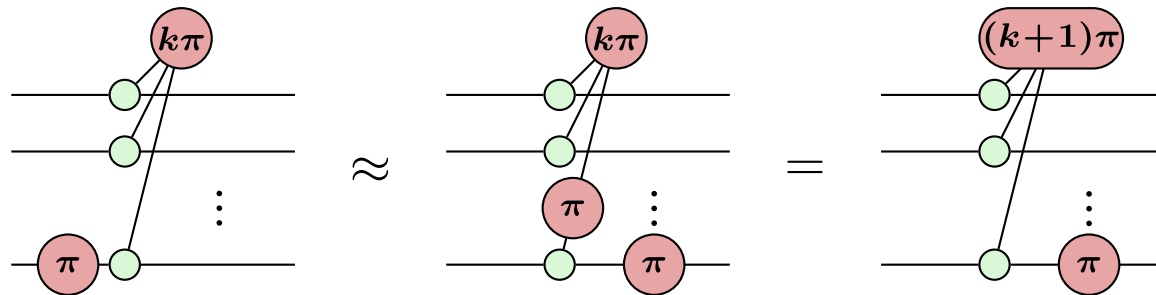
Modelling errors

Pauli errors are modelled by introducing X, Y, or Z flips on edges in a ZX-diagram:



Modelling errors

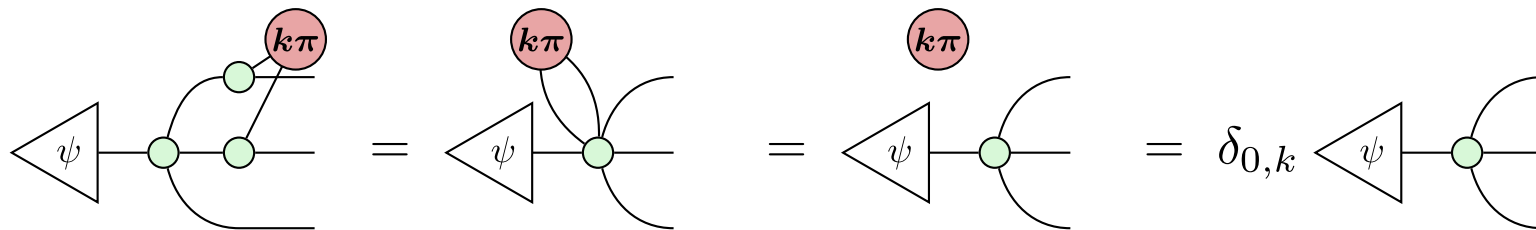
Anti-commuting errors flip measurement outcomes:



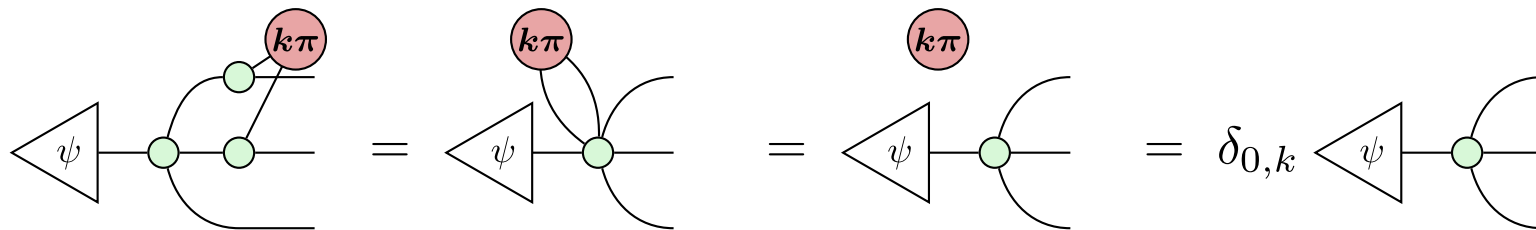
Example: GHZ code



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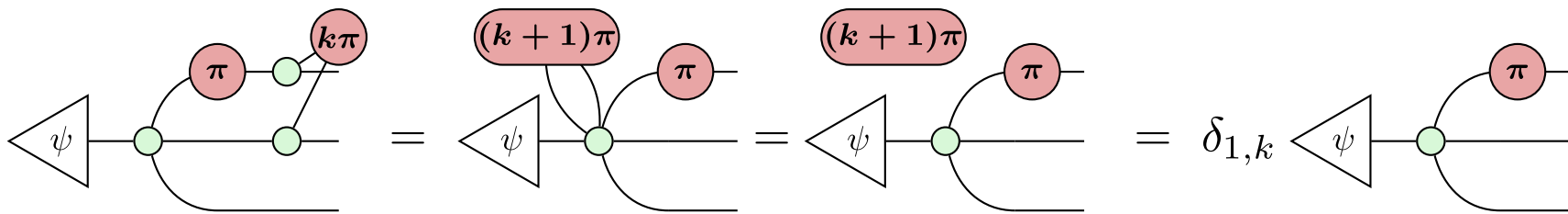


Example: GHZ code



This is a **stabiliser** measurement. In the absence of errors, it doesn't do anything!

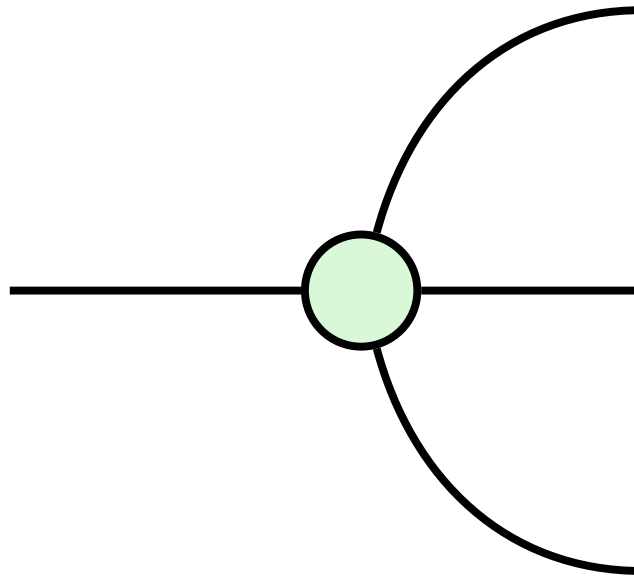
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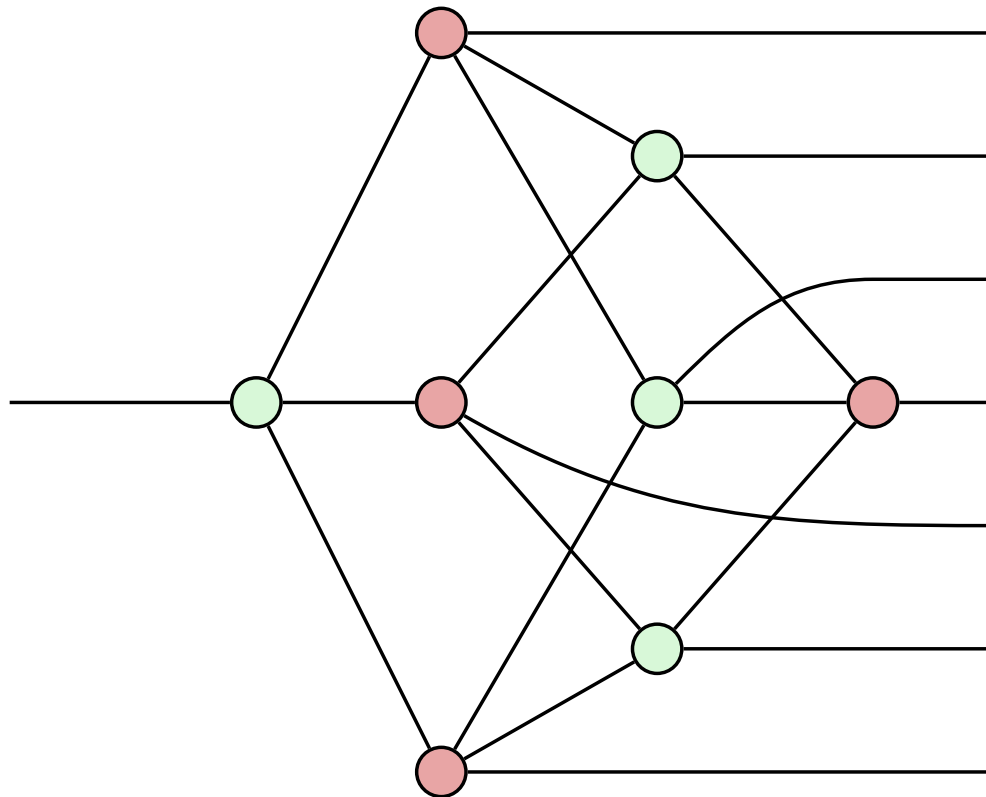
...but some errors will flip the measurement outcome,
giving an **error syndrome**.

Graphical encoders

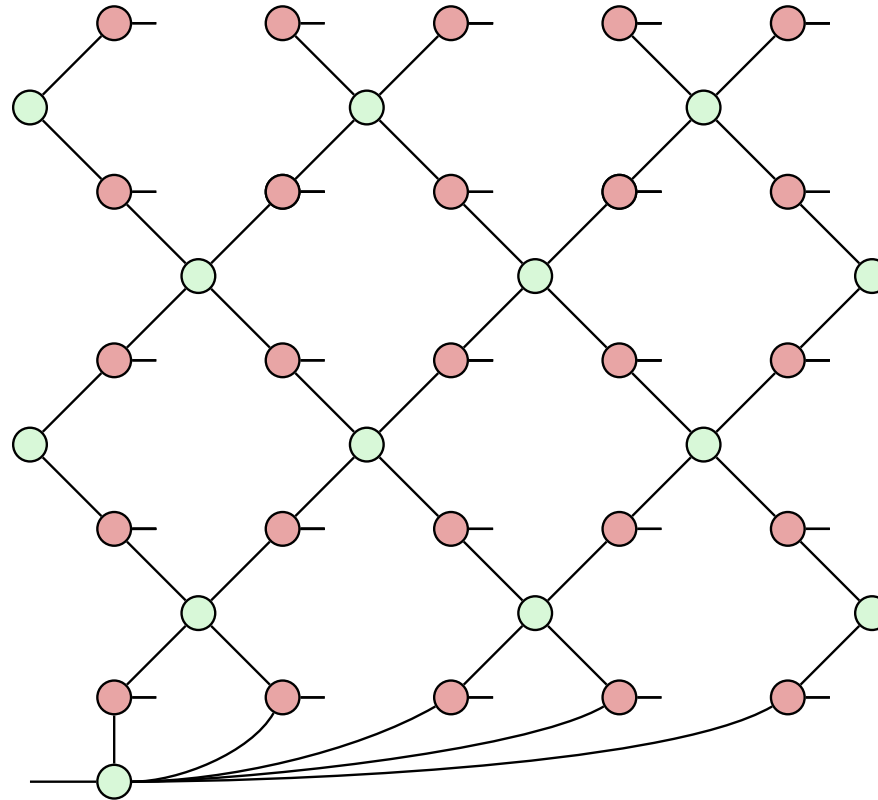
...give us a simple dictionary between QEC codes and ZX-diagrams, e.g. the GHZ code can be fully represented as:



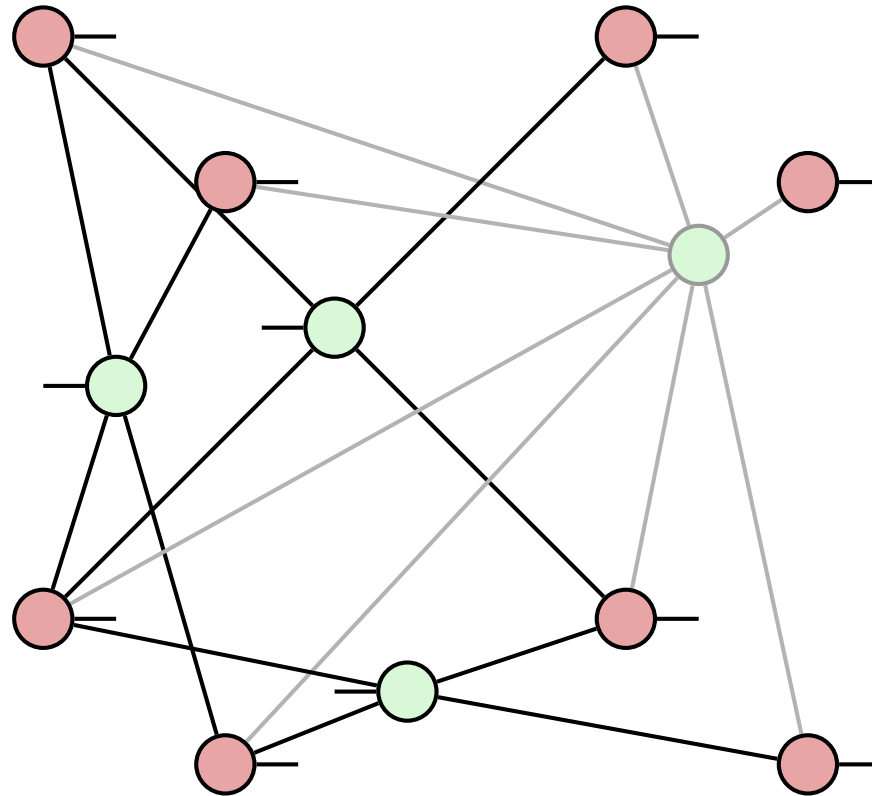
Example: Steane code



Example: Surface code



Example: $[[8, 3, 2]]$ Colour code



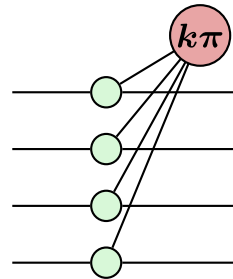
Fault-tolerant computation

But codes are only half the story in FTQC. We also need to know how to implement operations fault-tolerantly.

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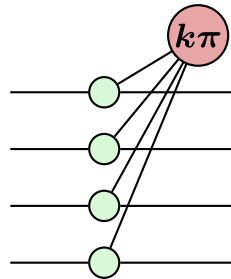
Consider a 4-qubit $Z \otimes Z \otimes Z \otimes Z$ measurement:



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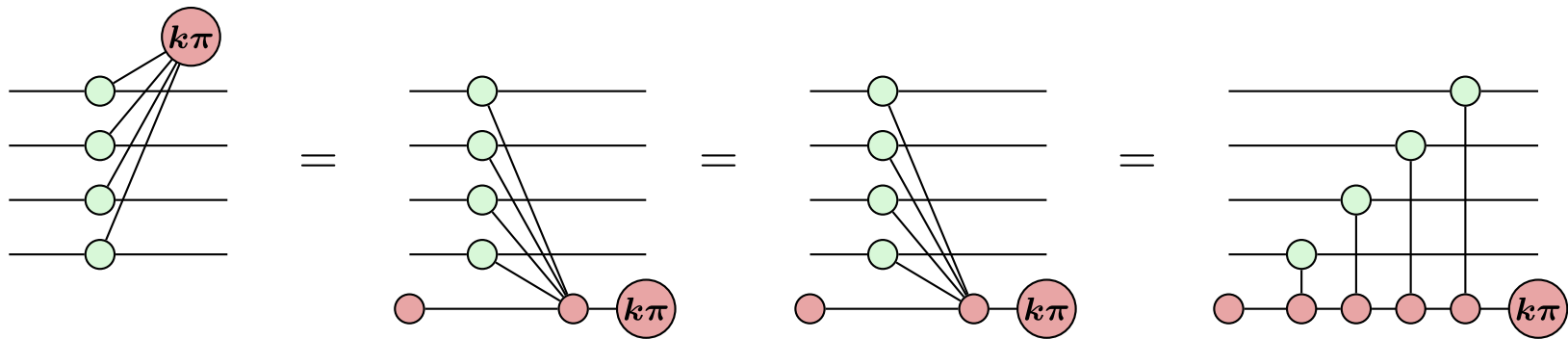
Consider a 4-qubit $Z \otimes Z \otimes Z \otimes Z$ measurement:



This isn't a basic operation (for most quantum computers).
How can we implement this?

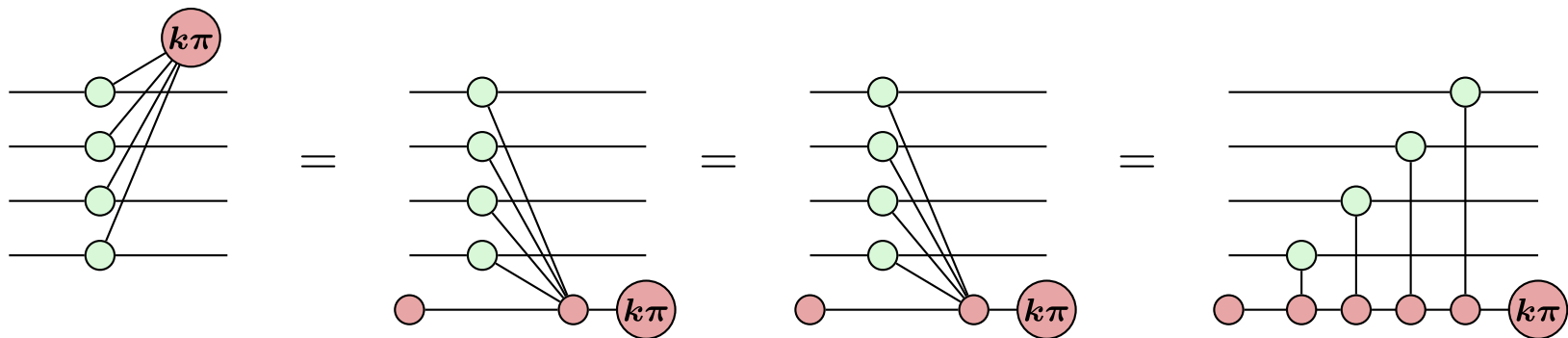
Example: measurement circuits

ZX can give us an answer:



Example: measurement circuits

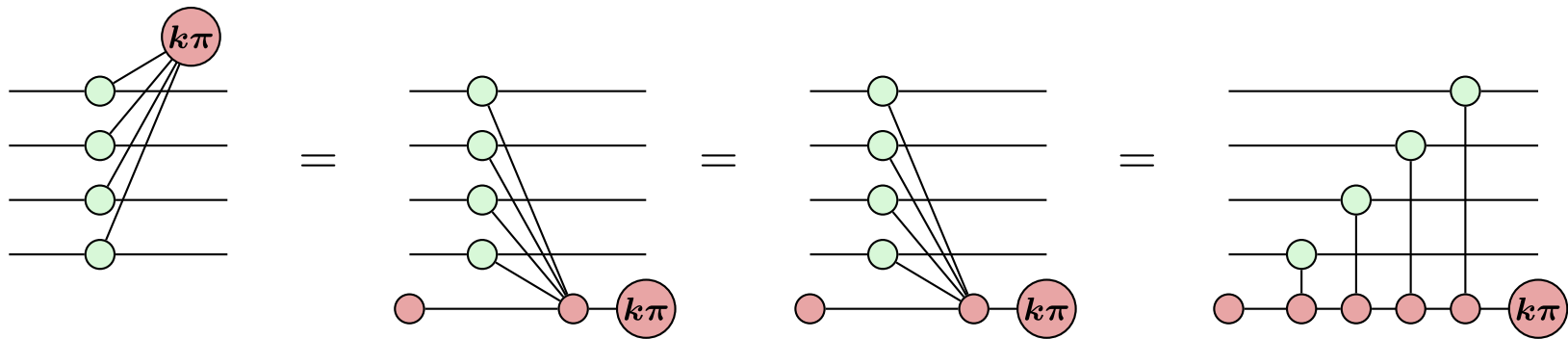
ZX can give us an answer:



Q: Is the LHS really equivalent to the RHS?

Example: measurement circuits

ZX can give us an answer:

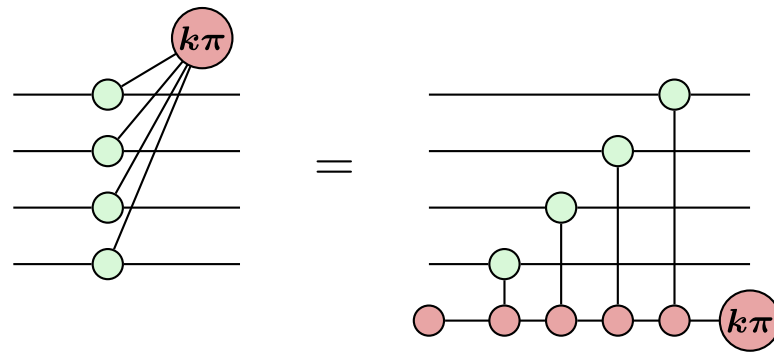


Q: Is the LHS really equivalent to the RHS?

A: It depends on what "equivalent" means.

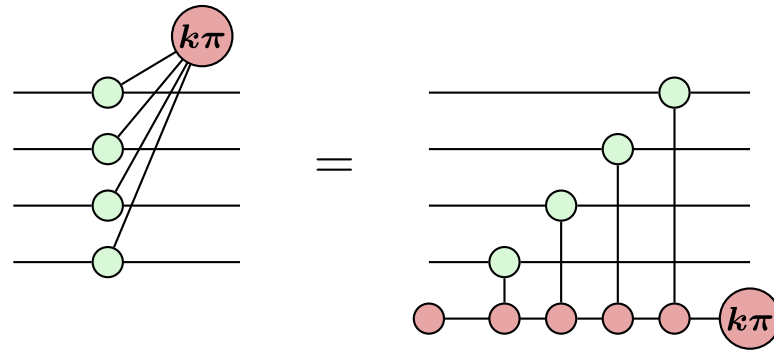
Example: measurement circuits

They both give the same linear map, i.e. they behave the same in the absence of errors.

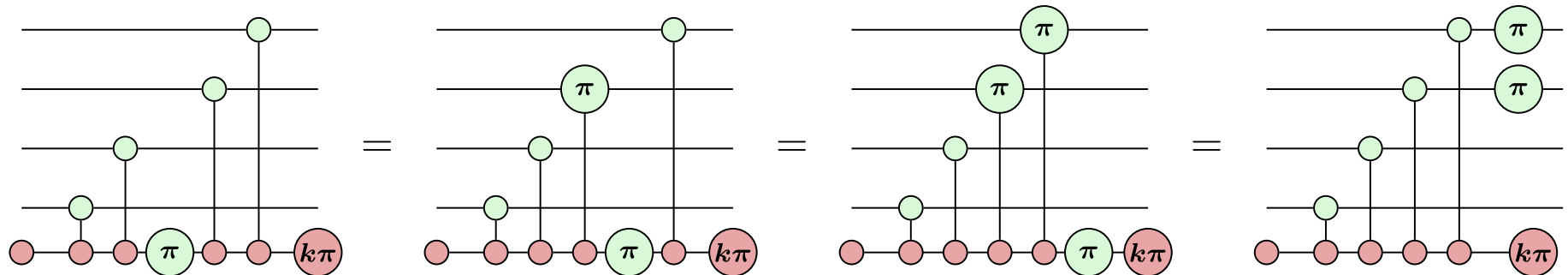


Example: measurement circuits

They both give the same linear map, i.e. they behave the same in the absence of errors.



But they behave differently in the **presence** of errors, e.g.



Solution: Fault equivalence

What we need is a notion of equivalence that captures the behaviour of circuits (or ZX-diagrams) in the presence of errors, **fault-equivalence**:

$$D \hat{=} E$$

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$$D \hat{=} E$$

This is a **finer-grained** notion of equivalence than the usual one:

$$D \hat{=} E \implies D = E$$

$$D = E \not\implies D \hat{=} E$$

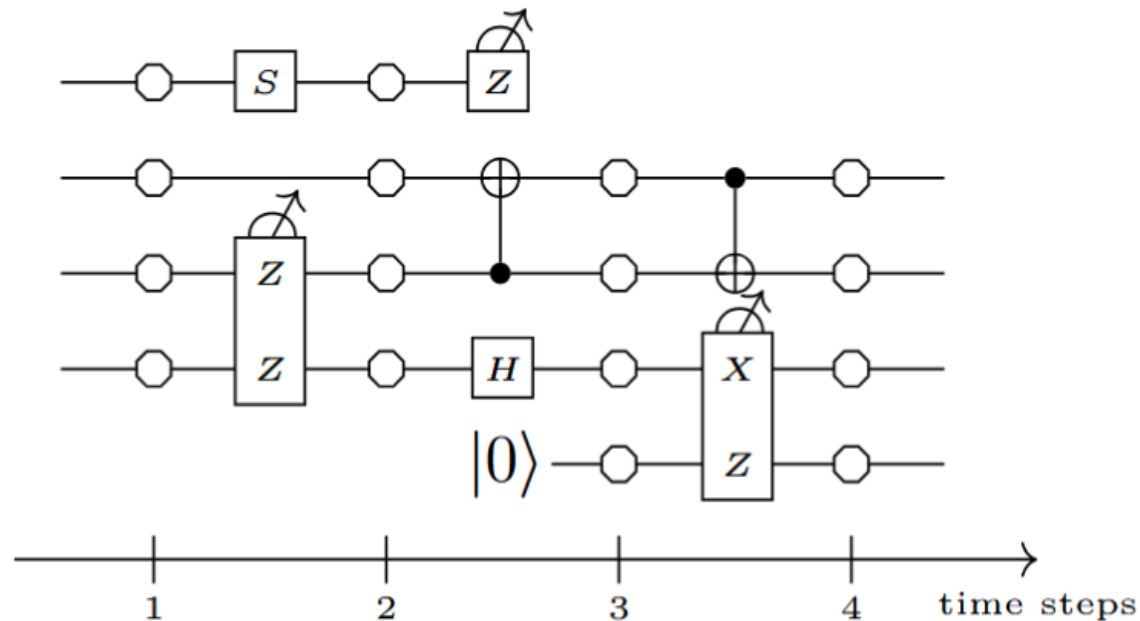
Error locations

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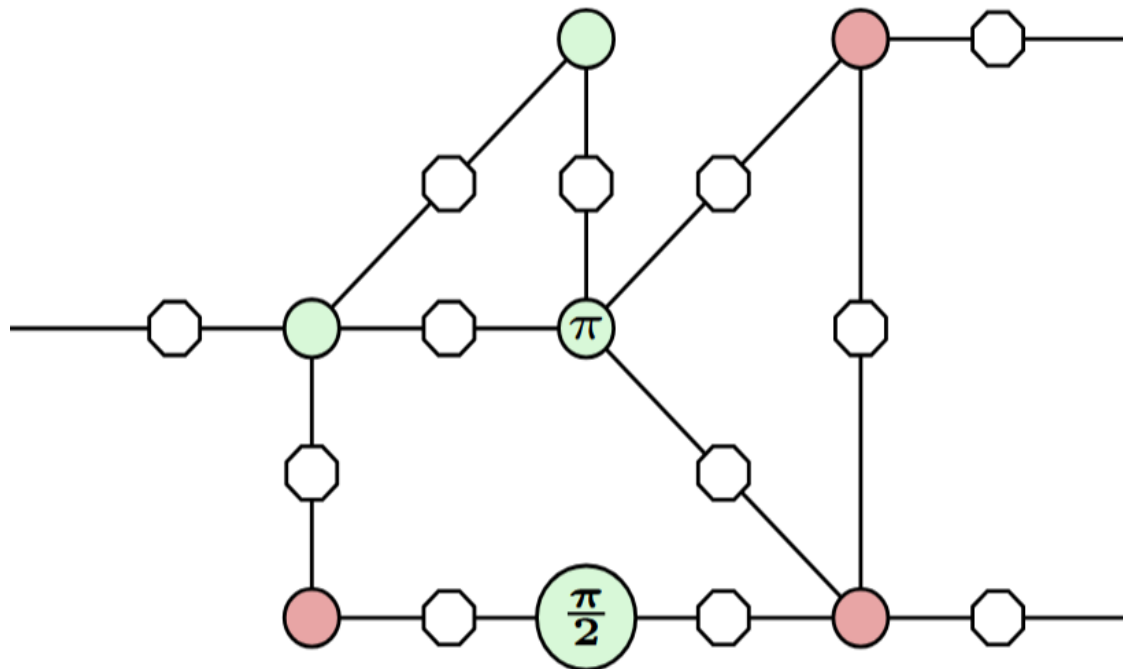
For circuits, fault locations are qubits $\times$ time steps:



Error locations

Similar to [space-time codes](#), we can model faults by tracking their locations in a circuit or ZX-diagram.

For ZX-diagrams, locations are time-agnostic, we allow them at any edge in the diagram:



Applying a Pauli to a Circuit/Diagram

- For a circuit/diagram with a set $\mathcal{L}$ of fault locations, a fault is a Pauli

$$F \in \mathcal{P}^{|\mathcal{L}|}$$

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- We can **apply** any Pauli $F \in \mathcal{P}^{|\mathcal{L}|}$ to a diagram D by plugging it in, to make a linear map $D[F]$
 - *convention: interpret measurements as projection onto $+1$, and assume $D[I] \neq 0$*

Applying a Pauli to a Circuit/Diagram

- For a circuit/diagram with a set $\mathcal{L}$ of fault locations, a fault is a Pauli

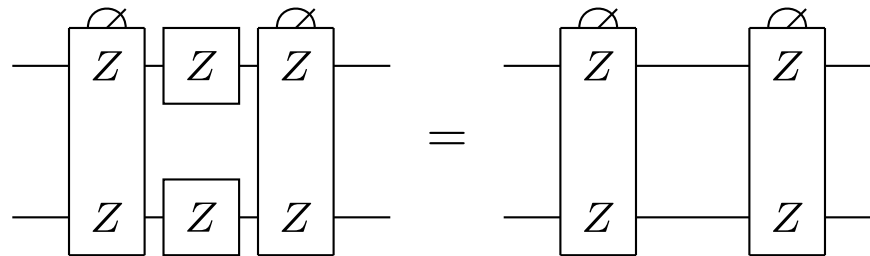
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 - *convention: interpret measurements as projection onto $+1$, and assume $D[I] \neq 0$*
- **There are 3 possibilities...**

F does nothing

$$D[F] = D[I]$$

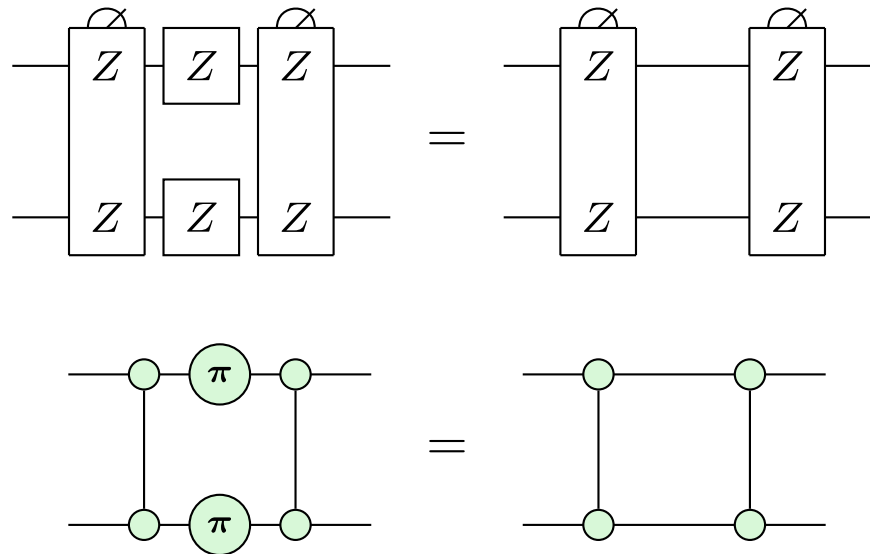
i.e. it is a **gauge** of the diagram



F does nothing

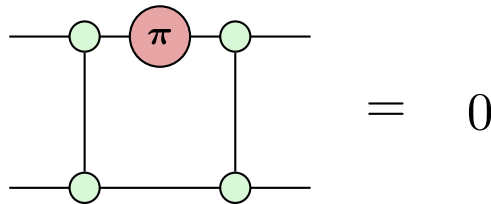
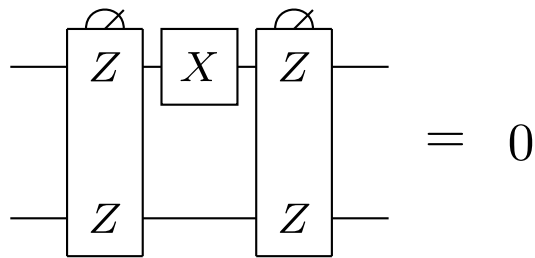
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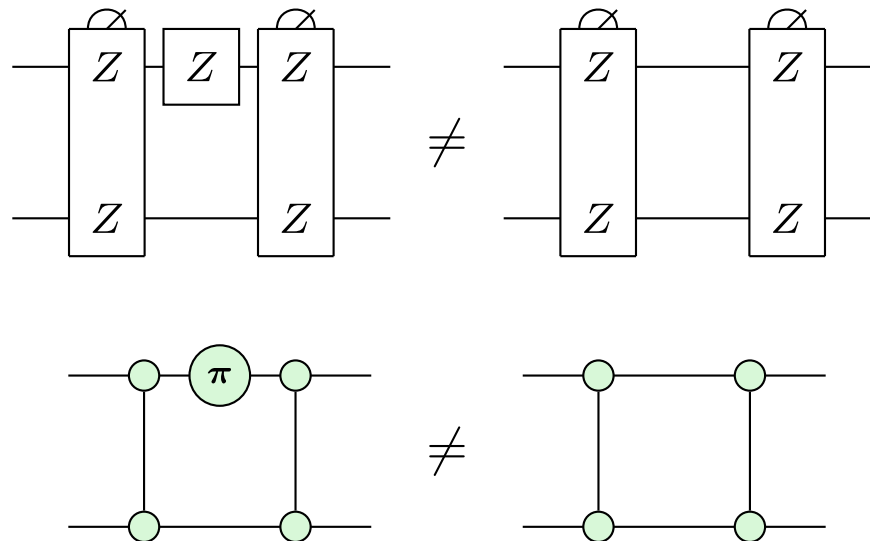
F is detectable

$$D[F] = 0$$



F is undetectable

$$D[F] \neq 0 \text{ and } D[F] \neq D[I]$$



Pauli error models

- A **Pauli error model** is a choice of basic faults $F_1, \dots, F_b$. The **weight** of a fault is the number of basic faults in its product.

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- A **Pauli error model** is a choice of basic faults $F_1, \dots, F_b$. The **weight** of a fault is the number of basic faults in its product.

- Ex: **Phenominological/naive model**:

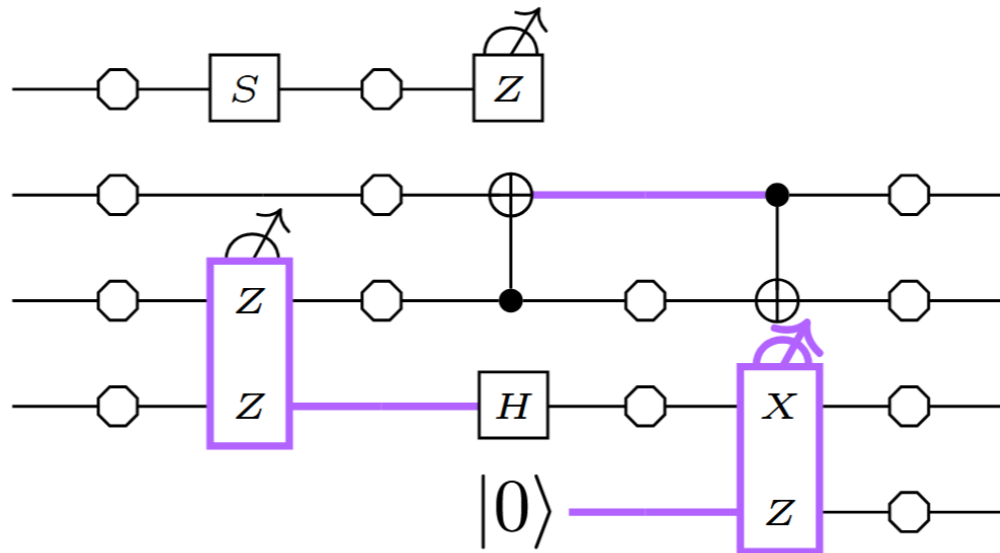
basic faults $:=$ all single-qubit Paulis

- Ex: **Circuit-level noise model**:

basic faults $:=$ all single-qubit Paulis (memory errors) +
some multi-qubit Paulis (correlated gate/measurement errors)

Pauli error models

Ex: **sub-models**, where we remove some basic faults to represent components we assume are noise-free:



Fault-equivalence

Definition: Two circuits (or ZX-diagrams) C, D are called *fault-equivalent*:

$$C \hat{=} D$$

if for any **undetectable** fault F of weight w on C , there exists an undetectable fault F' of weight $\leq w$ on D such that $C[F] = D[F']$ (and vice-versa).

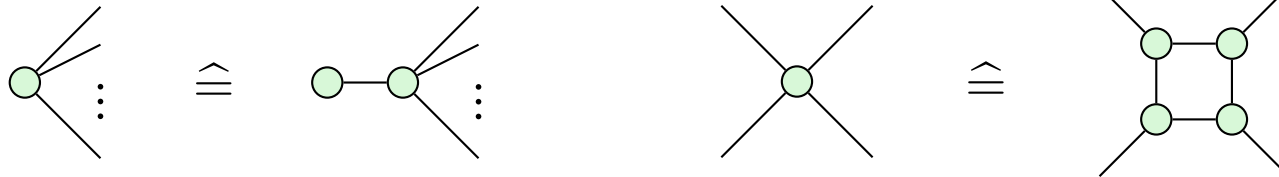
Fault-equivalence

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Fault-equivalence

- While all the ZX rules preserve map-equivalence, only some rules preserve fault-equivalence.
- It turns out the ones that do, e.g.

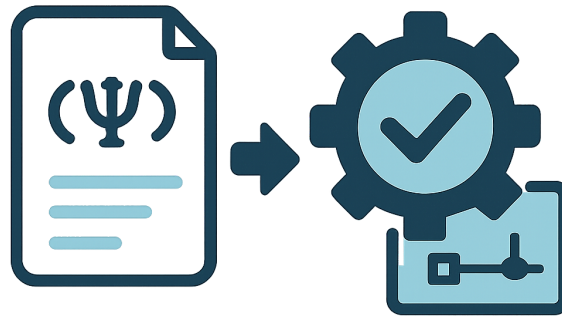


...are very useful for compiling fault-tolerant circuits!

Paradigm: Fault-tolerance by construction

[arXiv:2506.17181](https://arxiv.org/abs/2506.17181)

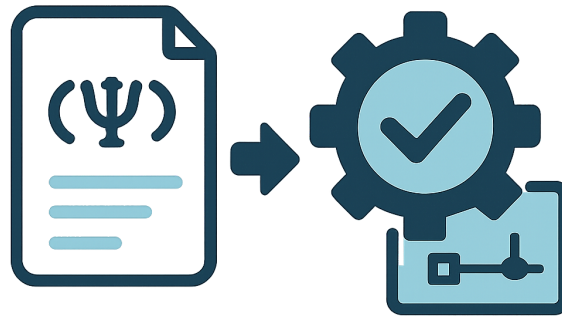
Idea: start with an idealised computation (i.e. specification) and refine it with fault-equivalent rewrites until it is implementable on hardware.



Paradigm: Fault-tolerance by construction

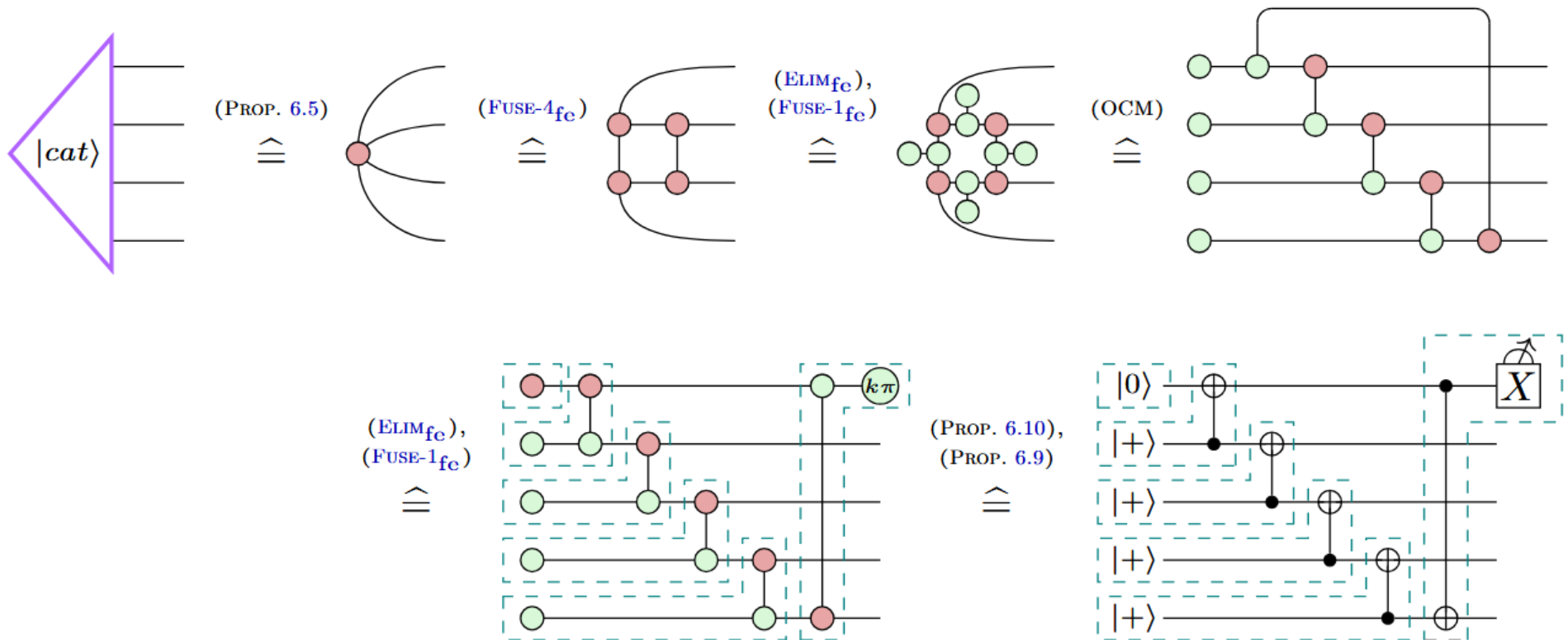
arXiv:2506.17181

Idea: start with an idealised computation (i.e. specification) and refine it with fault-equivalent rewrites until it is implementable on hardware.

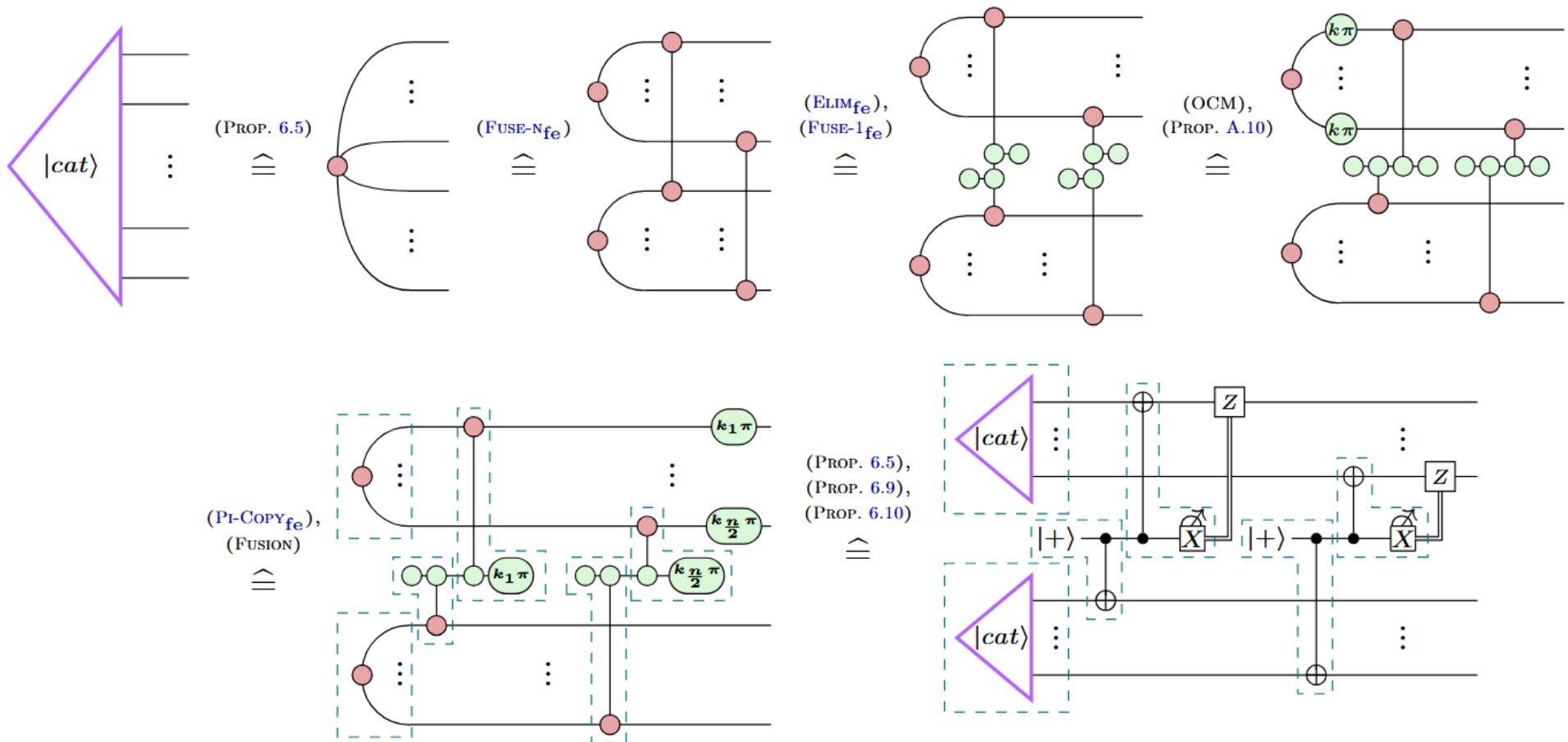


Specification/refinement has been used in formal methods for classical software dev since the 1970s. **Why not for FTQC?**

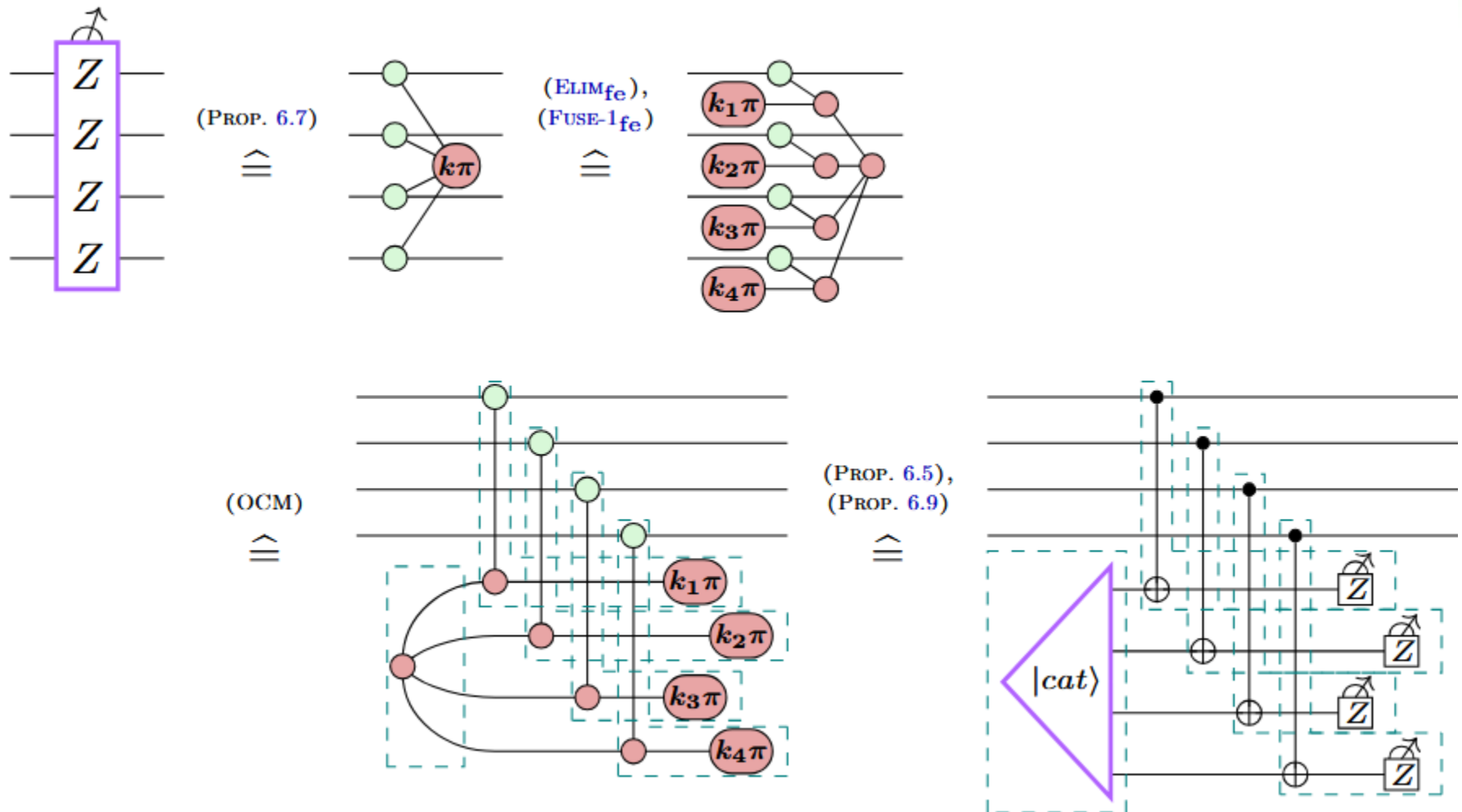
Example: Cat state preparation



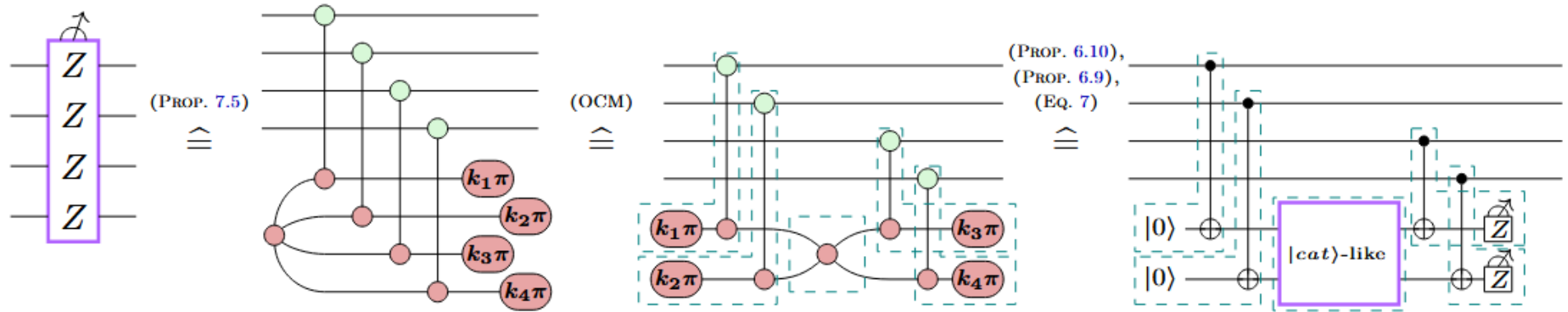
Example: Cat state preparation



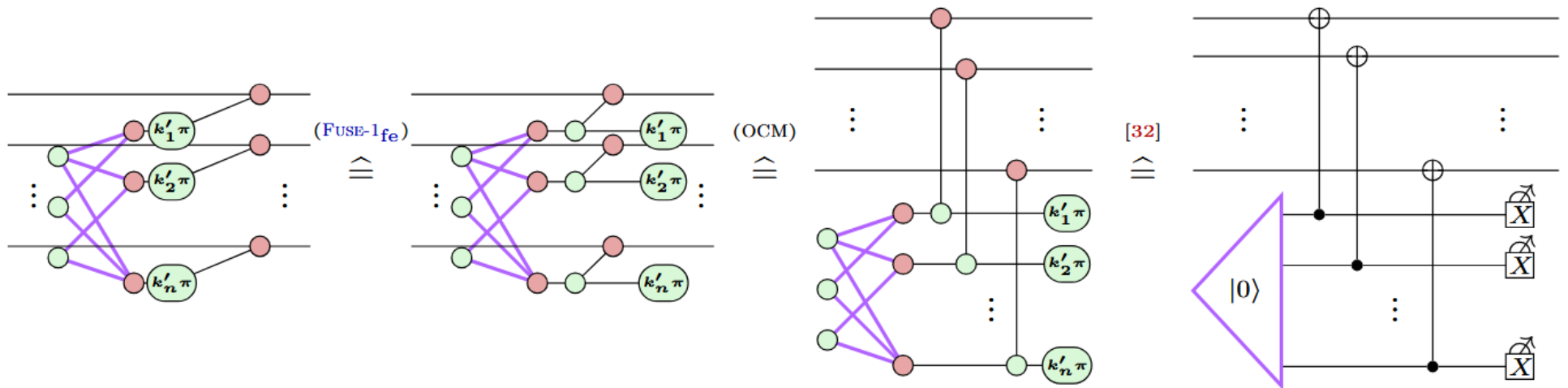
Example: Shor-style syndrome extraction



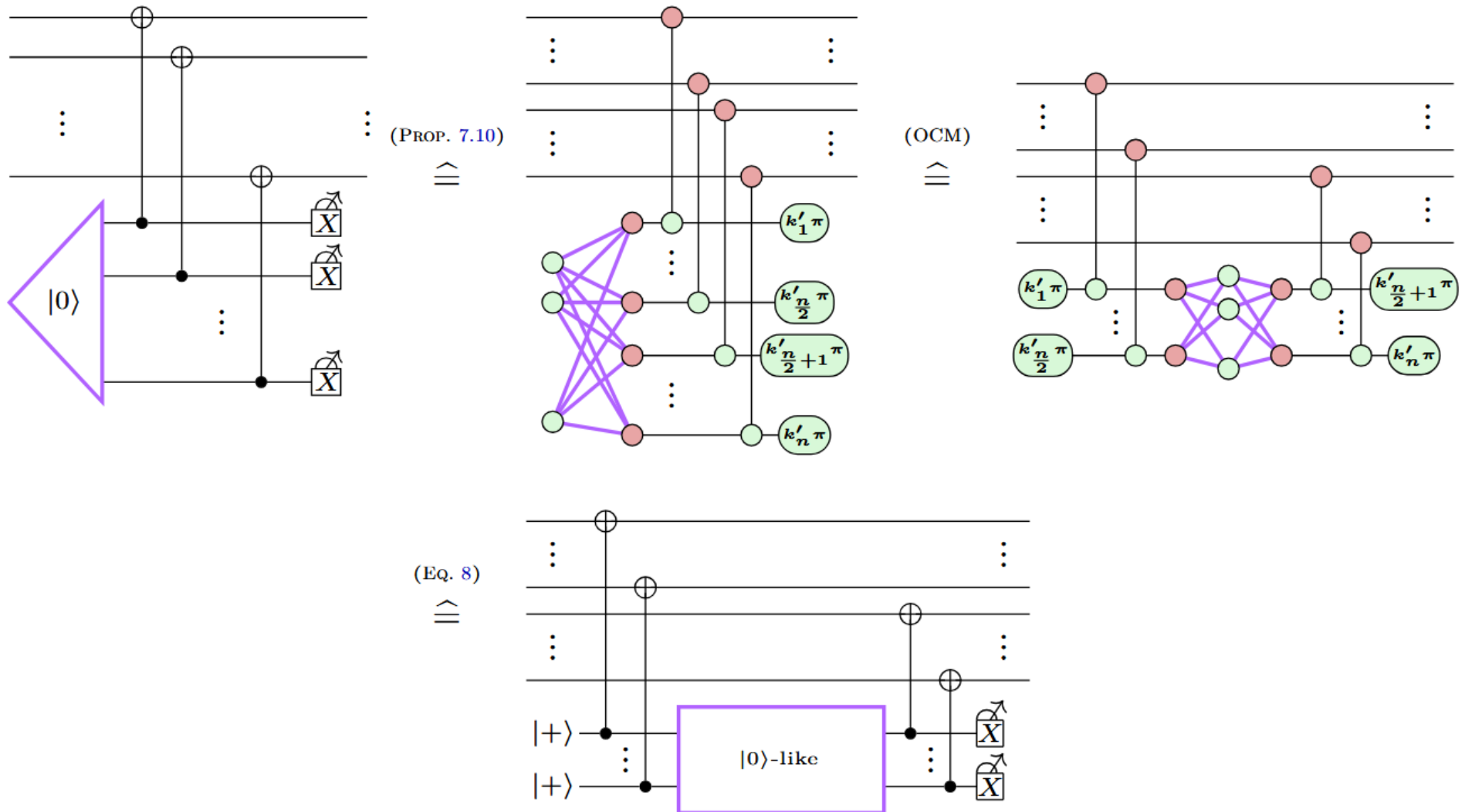
Example: A new variation on Shor



Example: Steane-style syndrome extraction



Example: A new variation on Steane (the same trick)

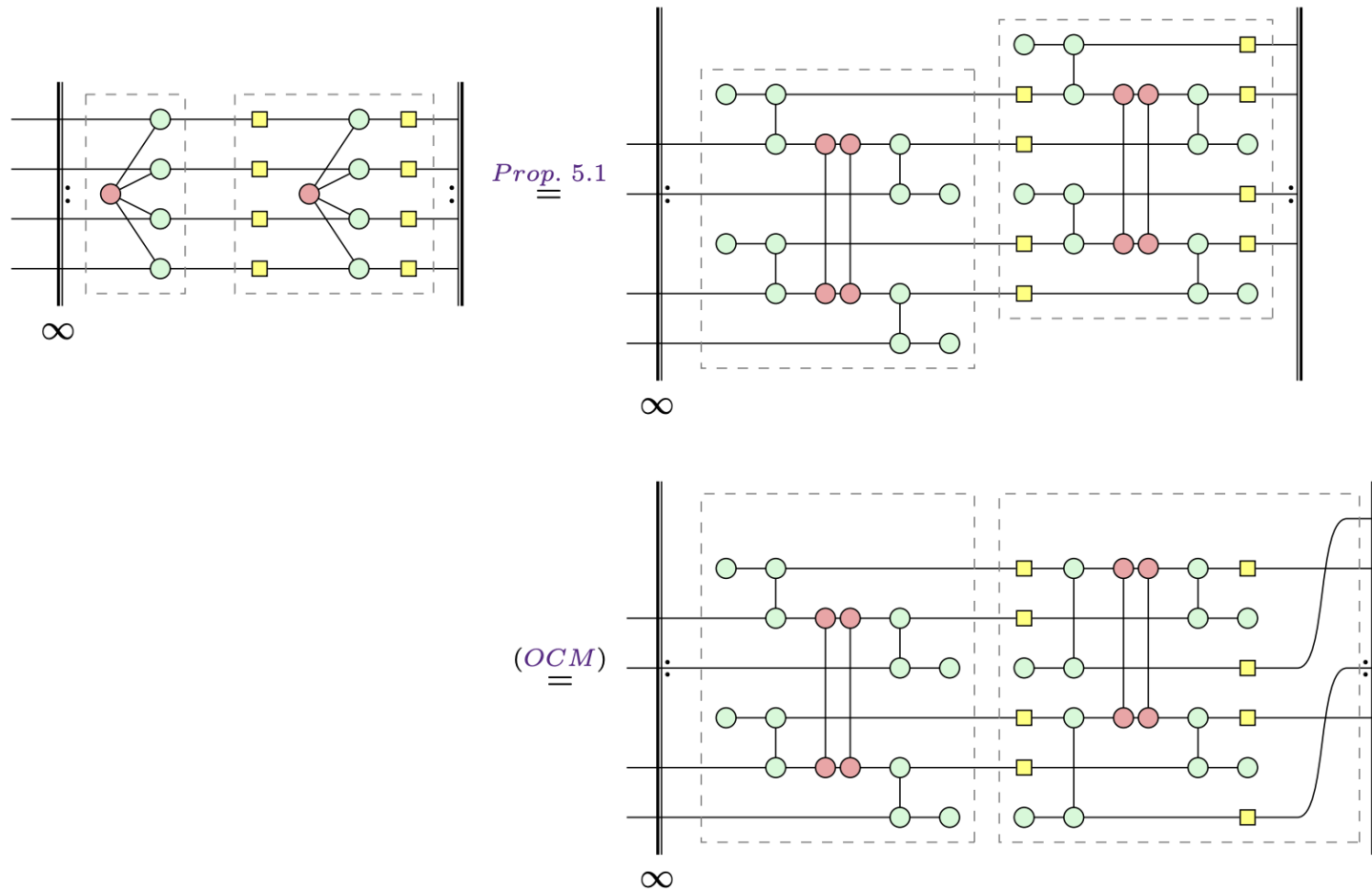


...with fewer qubits and a lower logical error rate!

FTbC Survey

Floquetification

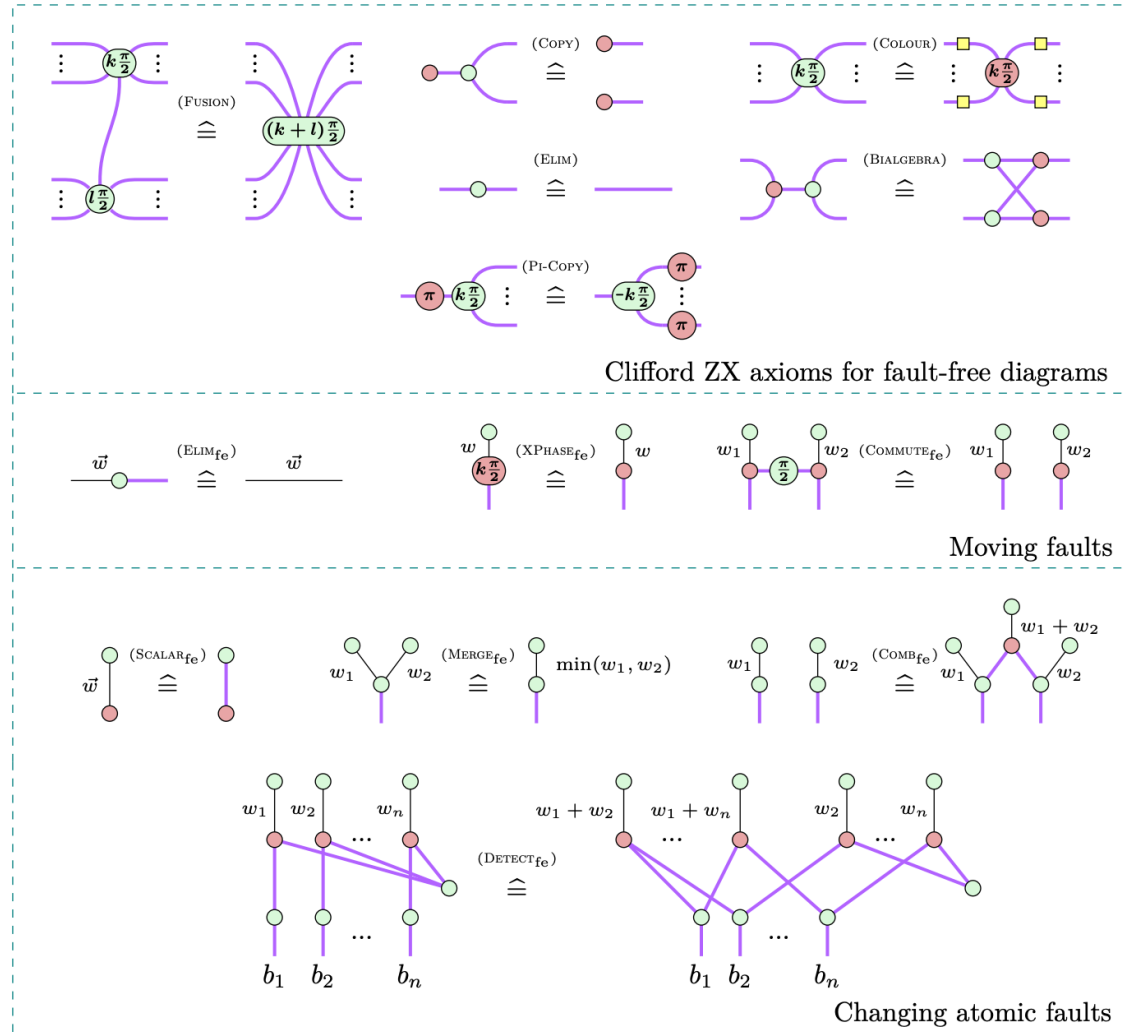
B Rodatz, B Poór, AK. [arXiv:2410.17240](https://arxiv.org/abs/2410.17240)



See also: A. Townsend-Teague, J. Magdalena de la Fuente, M. Kesselring. [arXiv:2307.11136](https://arxiv.org/abs/2307.11136)

Completeness for Fault-equivalent Rewriting

M Rüsç, AK, B Rodatz. [arXiv:2510.08477](https://arxiv.org/abs/2510.08477)

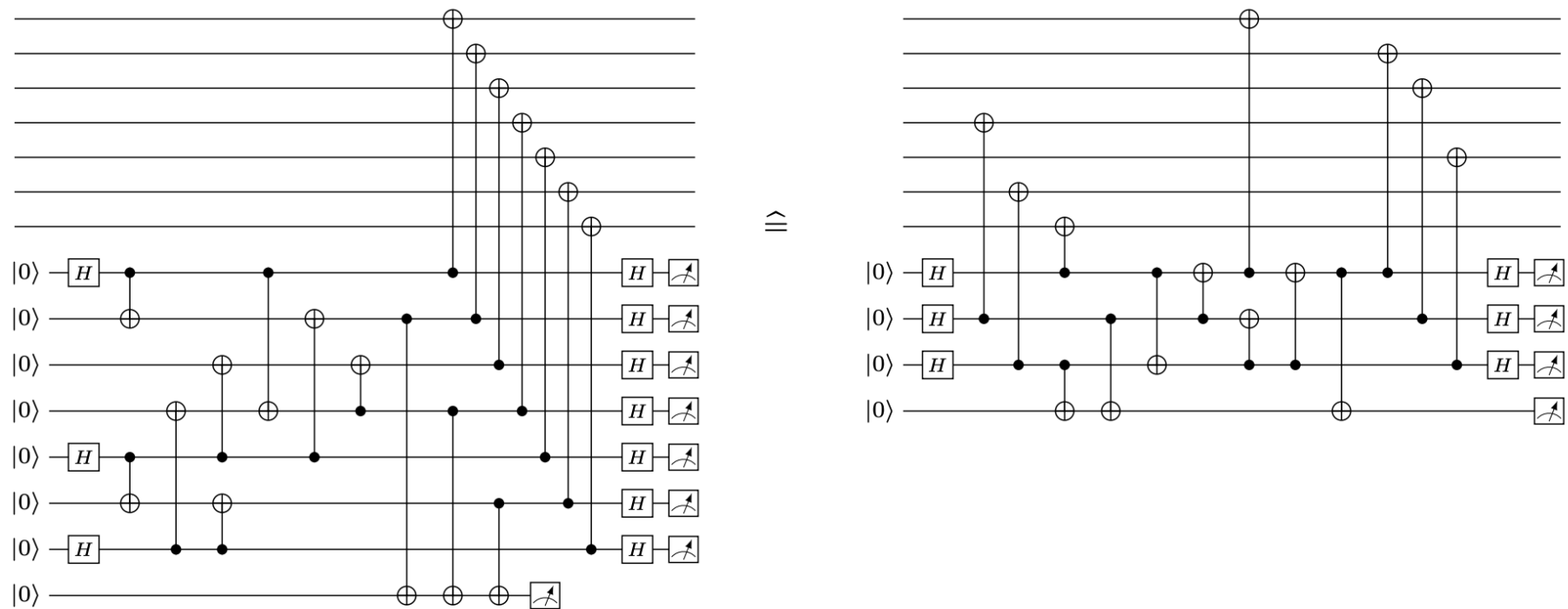


Ultra-low overhead syndrome extraction for the Steane Code

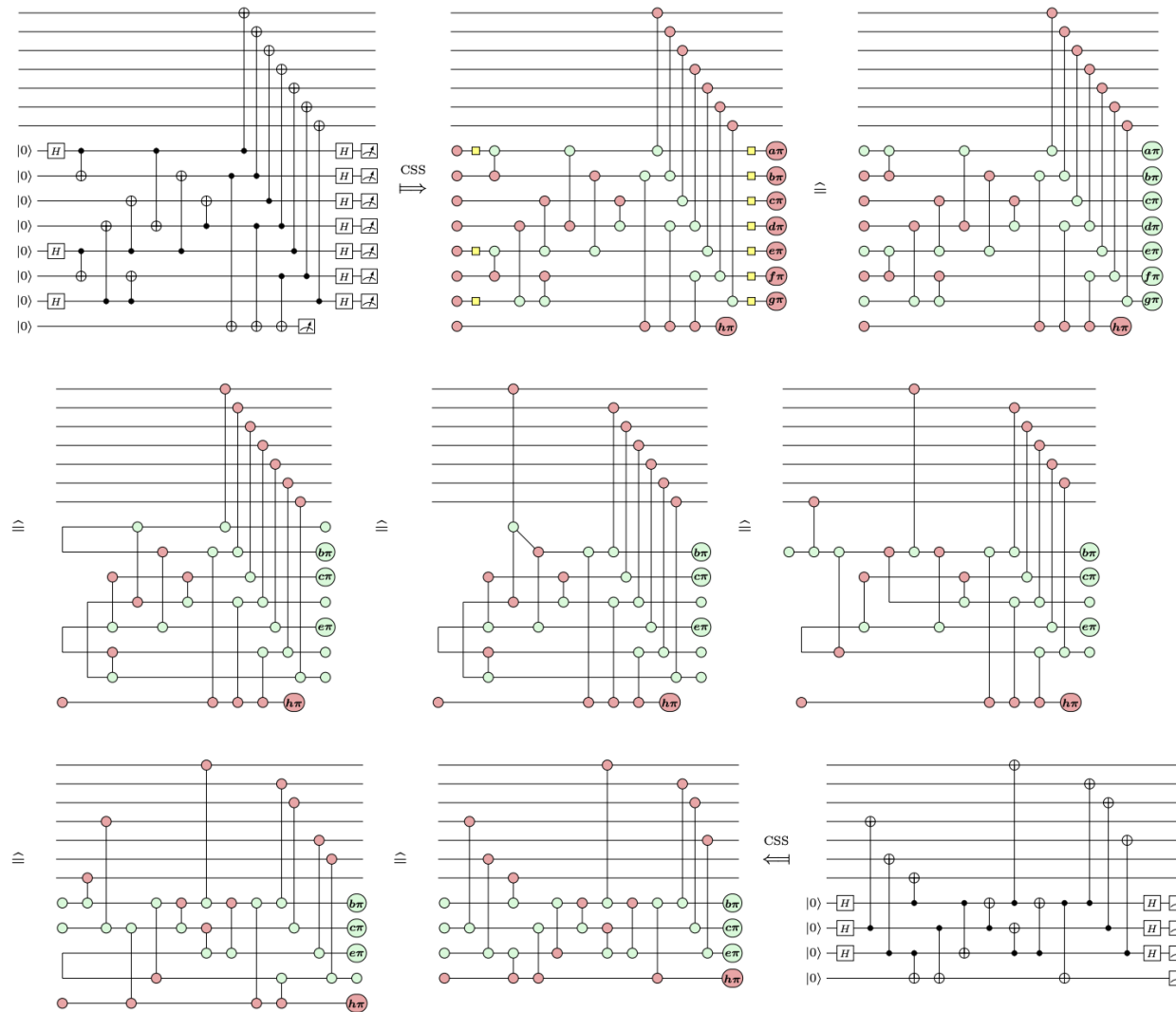
B Poór, B Rodatz, AK. [arXiv:2511.13700](https://arxiv.org/abs/2511.13700)

an 11-CNOT (optimal) syndrome extraction circuit for the Steane code*

Theorem 1. The following transformation is fault-equivalent:



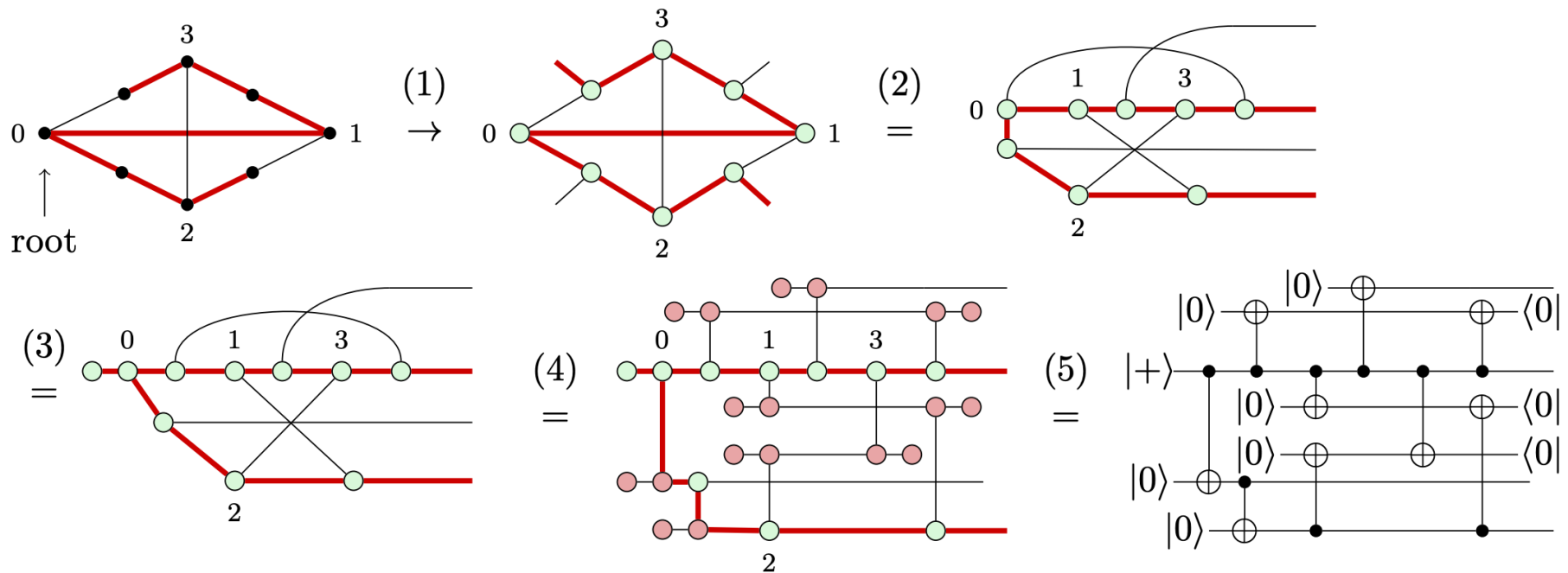
Ultra-low overhead syndrome extraction for the Steane Code



SpiderCat: Optimal FT Cat State Preparation

A Khesin, S Li, B Poór, B Rodatz, J van de Wetering, R Yeung.

[arXiv:2603.05391](https://arxiv.org/abs/2603.05391)



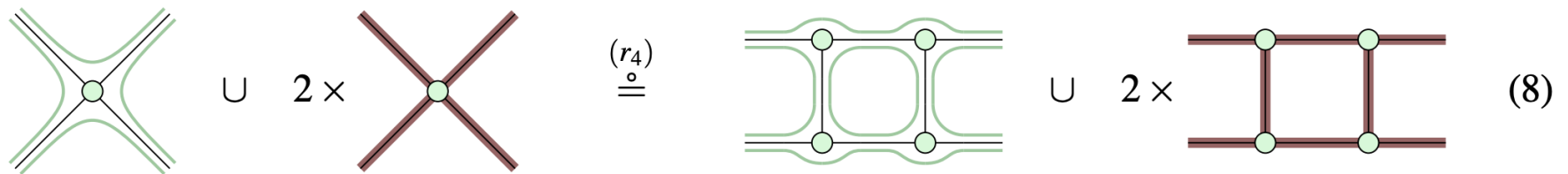
gives CNOT lower bounds for all sizes n and max FT error weight t

+ optimal explicit constructions for most pairs
 $(n \leq 100, t \leq 5)$ and $(n \leq 50, t \leq 7)$

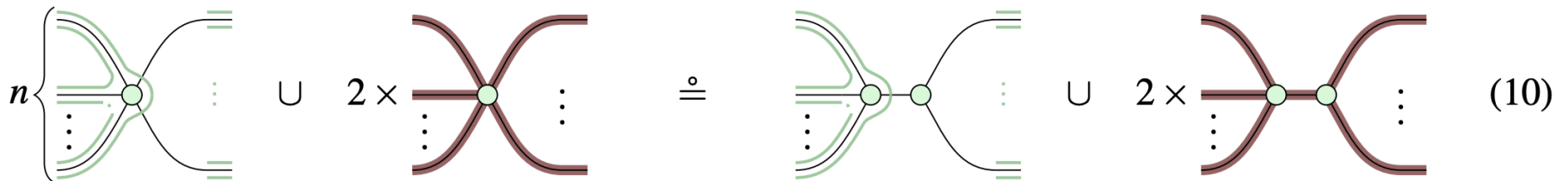
MWPM-Decodability Preserving Rewriting

M Schweikart, L Grans-Samuelsson, AK, B Rodatz. [arXiv:2603.19522](https://arxiv.org/abs/2603.19522)

Proposition 3.13 (r_4). *The following rewrite is CSS matchability preserving:*



Proposition 4.3. *Let $C(D)$ have a circuit distance of more than n . Then, the following rewrite is matchability-preserving:*

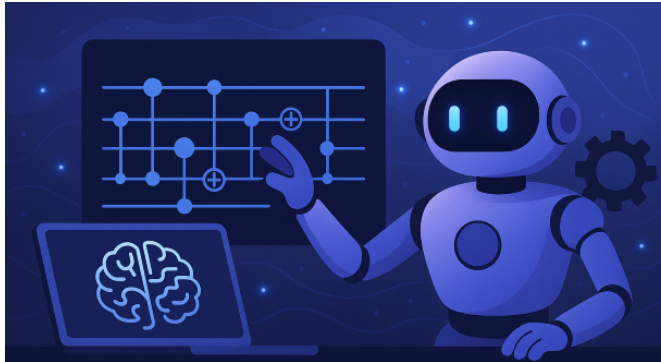


matchable QEC codes $\Rightarrow$ FT syndrome extraction circuits with matchable DEMs

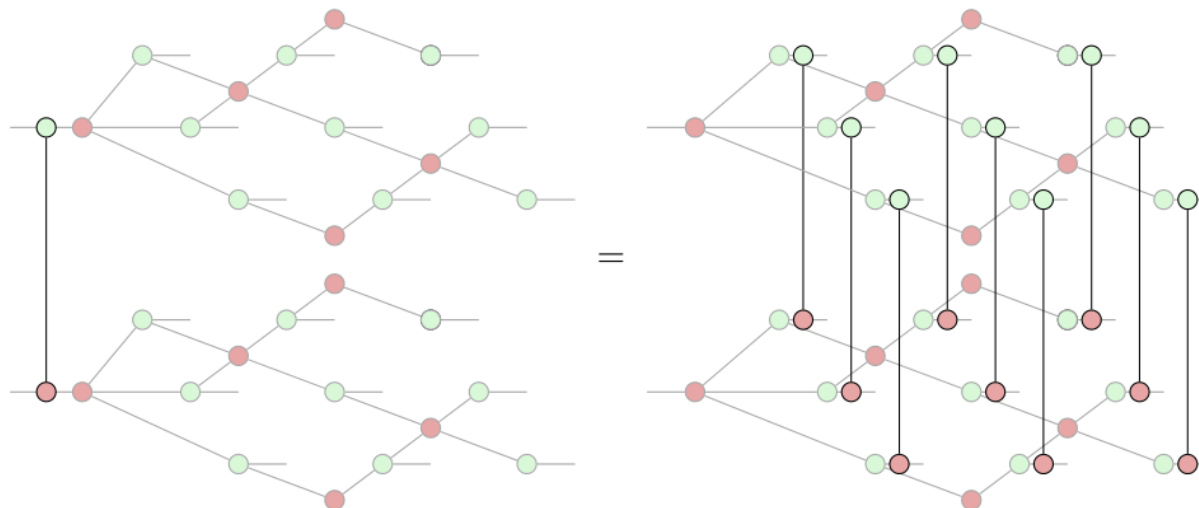
Lots to do!

Lots to do!

- Automatically building/optimising FT circuits (e.g. via heuristic search or AI)

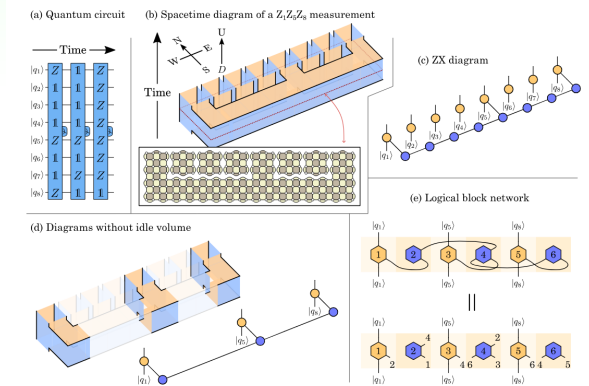
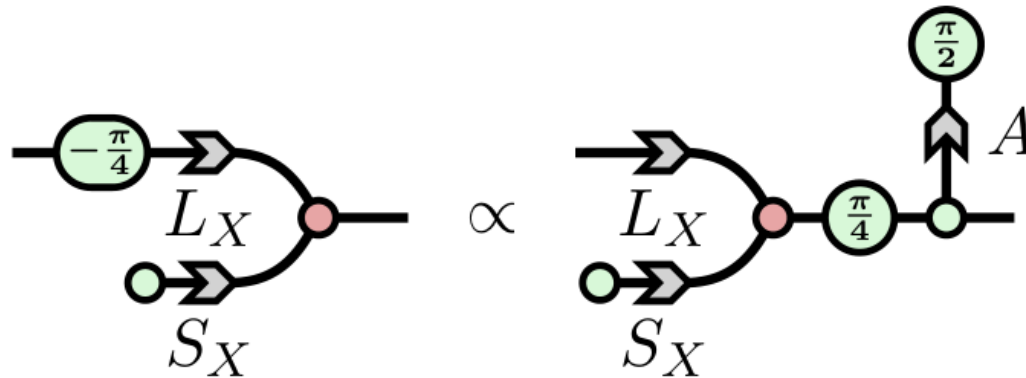


- Logical computation and measurement (esp. in non-traditional QEC paradigms, like dynamical codes)



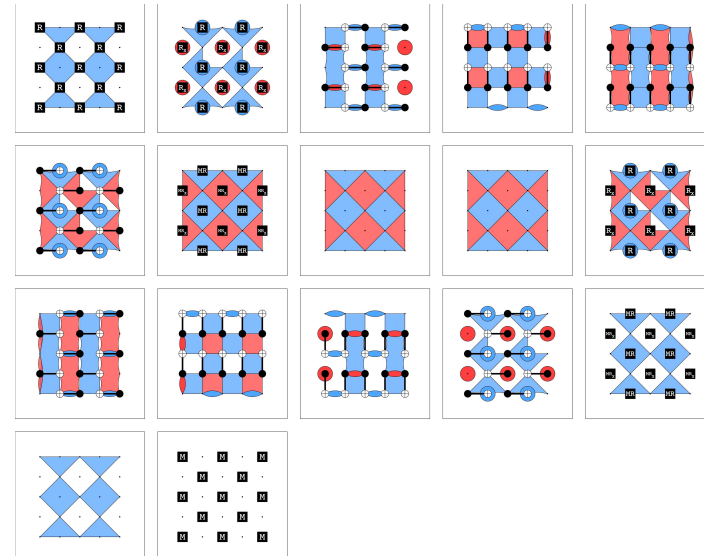
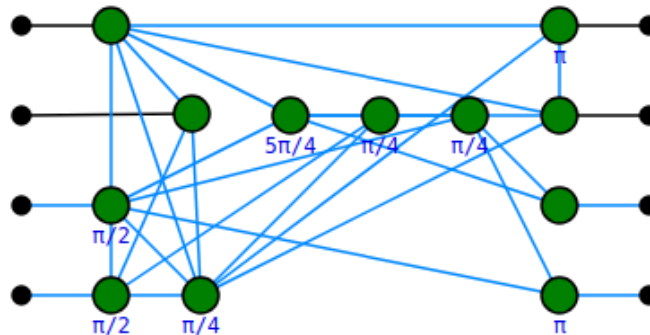
Lots to do!

- Scalability, compositionality, and FT architectures



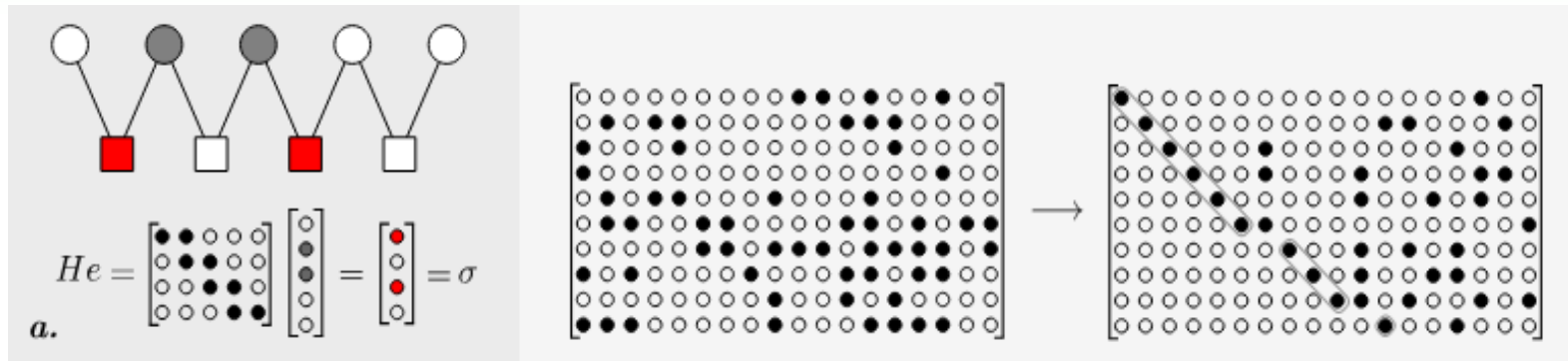
- Tooling, automation, and integration
(PyZX/QuiZX/ZXLive/Paritea/Stim)

```
In [11]: zx.simplify.clifford_simp(g)
          g.normalise()
          zx.d3.draw(g)
```

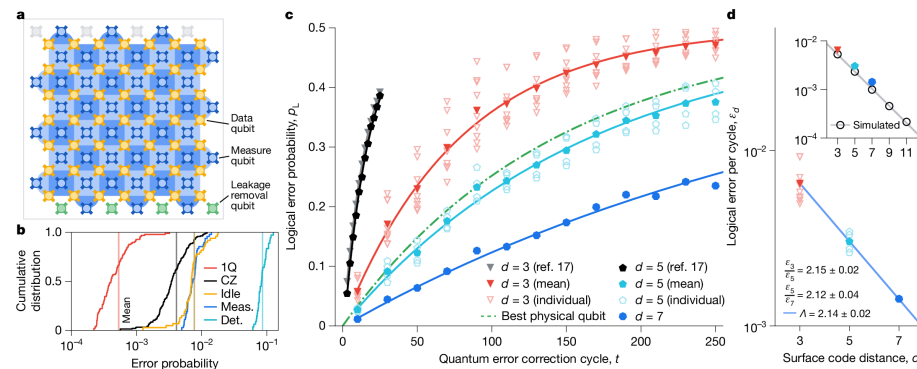


Lots to do!

- Decoding errors and preserving efficient decodability (beyond matchable codes/circuits)



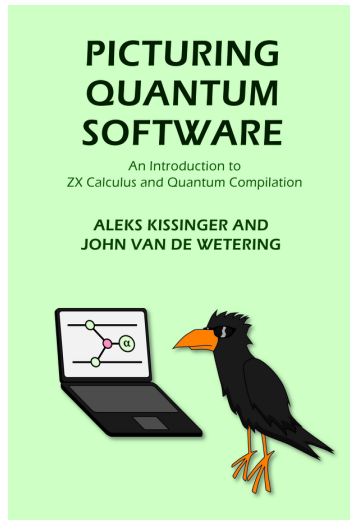
- Reasoning about stochastic noise (fault-equiv. = adversarial/worst-case noise behaviour)



Thanks!

Fault Tolerance by Construction » [arXiv:2506.17181](https://arxiv.org/abs/2506.17181)

Collabs: Andrey Khesin, Sarah Meng Li, Boldizár Poór, Benjamin Rodatz, John van de Wetering, Richie Yeung, Max Schweikart, Max Rüsçh, Linnea Grans-Samuelsson



<https://zxcalc.github.io/book>

(free book! Ch 12 = ZX + QEC)

<https://zxcalculus.com>

(350+ ZX papers, ~10% tagged QEC, online seminars, Discord)