

Combinatorial inverse folding

... when helix size matters, but energy models may be safely simplified?!

Yann Ponty

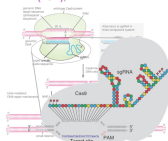
LIX, CNRS/Ecole Polytechnique

RiboNucleic Acids (RNA) in Human biology/health: Friend **and** Foe!

Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassaemia, Leber congenital amaurosis (LCA), cancers...



Hendel et al., 2015; Agrotis & Ketteler, 2015

Rational design

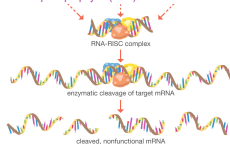
Regulation of gene expression

RNAi therapies (FDA approved)

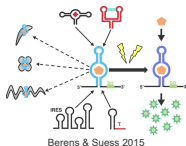
Primary hyperoxaluria type 1 (PH1),

Hereditary transthyretin amyloidosis (ATTRv),

Acute hepatic porphyria (AHP)



Encyclopaedia Britannica, Inc 2013



Sensor of metabolites
Riboswitches

RiboNucleic Acids (RNAs)



Encodes proteins

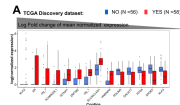
mRNA Vaccines

COVID-19, Malaria (Zika, CMV, Cancers?)

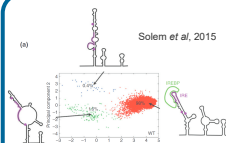
Quantitative expression

Transcriptomic signatures

Cancer diagnosis/prognosis/relapse...



[NGuyen et al, 2021]



Non-coding mutations

lncRNAs, miRNAs, structure-associated (RiboSnitches)

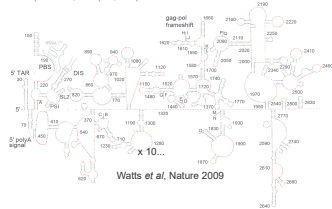
β -thalassaemia, duchenne muscular dystrophy,

Cystic fibrosis, Rett syndrome...

(2D) Structure Modeling

Genomic material for Human pathogens

HIV-1, SARS-CoV 2, HCoV, MERS



Definition (INVERSE-FOLDING(E) problem)

Input: Secondary structure R + Energy distance $\Delta > 0$

Output: RNA sequence $\omega \in \Sigma^*$ such that:

$$\forall R' \neq R \text{ compatible with } \omega : E(\omega, R') \geq E(\omega, R) + \Delta$$

or $\emptyset$ if no such sequence exists.

Difficult problem:

- ▶ Introduced in the 90s by Vienna/Leipzig wunderkids [Hofacker *et al* Monatshefte für Chemie 1994]
- ▶ NP-hardness in maximalist setting (Energy model as input [Schnall-Levin *et al*, ICML'08])
- ▶ NP-hardness for BP maximization with partial assignment [Bonnet *et al*, RECOMB'18]
- ▶ **Hard to study:** Non local, no theoretical framework, too many parameters. . .

Foreword and caveats

- ▶ *"All models are wrong, but some are useful"* George Box

Reductive models enable analysis. while hopefully retaining relevance

- ▶ In ViennaRNA we trust: RNAFold, not Nature (yet), scores

Ideally, uniform sampling from the preimage

- ▶ Is inverse RNA design *actually hard*? Worthy of further attention?

Molecular biology, Structural biology		→	YES
Comparative genomics, Computer Science. . .			
DNA computing		→	Meh...

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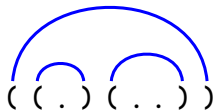
Inverse Folding in Base Pair maximization (maxBP) model

Definition (Inverse folding in maxBP)

Input: Target secondary structure R , i.e. set of pairwise non-crossing BPs

Output: RNA sequence $\omega \in \Sigma^*$ such that:

- ▶ Target R compatible with ω ;
- ▶ $\forall R' \neq R$, compatible with ω : $|R'| < |R|$.



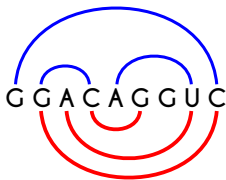
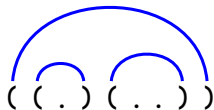
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Designability in simple BP-based energy models

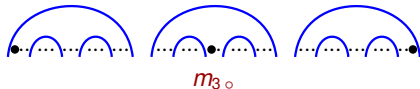
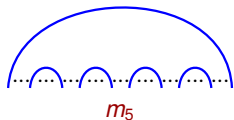
Partial characterization of **designable** structures [Hales *et al*, CPM'15 & Algorithmica'17]

- ▶ **Saturated structures (all positions paired):** Designable $\Leftrightarrow$ Multiloops degrees ≤ 4 (+ $\Theta(n)$ algo.)
- ▶ Designable $\Rightarrow$ Avoid multiloops with *degree* 5^+ (m_5), or *degree* 3^+ with 1^+ *unpaired* ($m_3 \circ$).

Designability in simple BP-based energy models

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- ▶ **Saturated structures (all positions paired)**: Designable $\Leftrightarrow$ Multiloops degrees ≤ 4 (+ $\Theta(n)$ algo.)
- ▶ Designable $\Rightarrow$ Avoid multiloops with **degree 5^+** (m_5), or **degree 3^+ with 1^+ unpaired** ($m_3 \circ$).



Remark: Non-designable motifs (obstructions) exist in all **energy model/design criterion**

Corollary [Yao *et al*, ACM-BCB'19&J Math Biol 2026]]

The fraction of designable structures is **exponentially decreasing** with length n

...**but** (fortunately) $\exists$ **infinite family** of **unsaturated designable** 2D structures [Jedwab *et al*, TCS 2020]

Separated coloring of RNA tree representations

Proper coloring of base pairs/nodes:

- ▶ Each base pair $\rightarrow$ one out of 3 colors: $\bullet \rightarrow G \cdot C$; $\circ \rightarrow C \cdot G$; $\bullet \rightarrow A \cdot U$ or $U \cdot A$.
- ▶ Parent node colored $\bullet$ (resp. $\circ$) cannot have $\circ$ (resp. $\bullet$) children;
- ▶ Children satisfy $\#\bullet \leq 1$, $\#\circ \leq 1$, $\#\bullet \leq 2$ and $\#\bullet + \#\circ < 2$

Definitions:

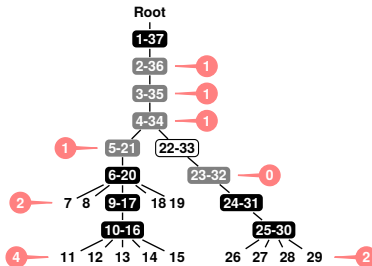
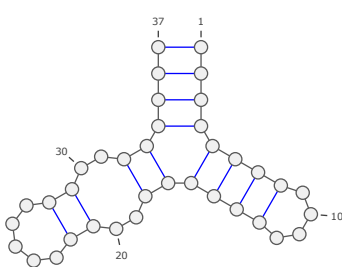
- ▶ **Level** of a base pair = $\#\bullet - \#\circ$ on path to root
- ▶ Coloring **separated** if proper and $\bullet$ base pairs and unpaired positions at **different** levels

Theorem ([Hales *et al.* Algorithmica 2017])

$\exists$ **Separated (proper)** coloring for structure $\Rightarrow$ Structure is designable in $\Theta(n)$ time

Separated Coloring

(((((((((.....)))))) ((.. ((.....)) ..))))))



GAAAAGUUGGUUUUCCUUCUCAGGUUUUCCUGUUUC

Intuition: All unpaired $\rightarrow$ A, so alternative structures need to pair every G, C, and A:

- If $\bullet$ base pairs conform to target, then $\bullet/\circ$ behave as saturated structure;
- If some $\bullet$ misbehaves, then some $\bullet$ must interact with leaf, delimiting regions with $\#G \neq \#C$, so we lose at least one G – C pair.

Theorem ([Boury *et al.* Alg Mol Biol 2025])

However, deciding if a structure is **separable** is **NP-hard**

Modulo m -separated coloring

Reminder: Separated coloring = induced **levels** of $\bullet$ and leaves do not overlap.

Idea: Prescribe **allowed values** for the **levels** mod m of $\bullet$ (set ξ) and leaves (set $\bar{\xi} := [0, m-1] \setminus \xi$).

Definition: Given m and $\xi \subseteq [0, m-1]$, a $(\xi, \bar{\xi})$ -separated coloring χ is such that, for each node v :

$$\blacktriangleright \chi(v) = \bullet \rightarrow \text{Level}_\chi(v) \bmod m \in \xi \quad \text{and} \quad \blacktriangleright v \text{ is leaf} \rightarrow \text{Level}_\chi(v) \bmod m \in \bar{\xi}.$$

Proposition (Boury *et al*, Alg Mol Biol 2026)

Given m and $(\xi, \bar{\xi})$, a $(\xi, \bar{\xi})$ -separated coloring can be found in $\Theta(m.n)$ -time using dyn. prog.

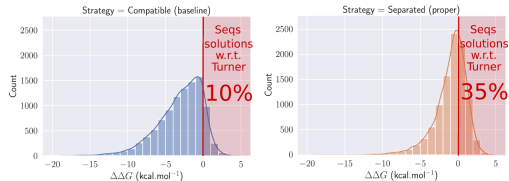
$\rightarrow$ (modulo) **m -separated colorings** can be found in $\Theta(2^m.n)$ time/ $\Theta(m.n)$ memory (FPT)

Also, any target structure with 3^+ BP/helix (+ many with helix length 2) is (modulo) 2-separated

Theorem (Boury *et al*, Alg Mol Biol 2026)

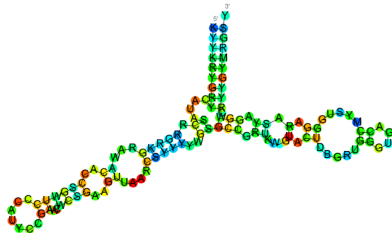
For targets w/o isolated stack/BP, maxBP inverse folding solvable in $\Theta(n)$ time

The good, the bad and the ugly



- Separated designs: **promising first step** towards inverse folding in **Turner** model
- Amenable to **uniform random generation**
→ Cardinality estimates for neutral network
- Good **seeding strategy** for local search
[Boury, . . . , Yao, RECOMB'25]

- maxBP considers **unrealistic competing structures** (e.g. isolated BPs)...
- ...and overlooks **real Turner competitors**
- Simple BP model **artificially restricts** loops
→ Crucial RNA structures are maxBP undesignable



Inverse folding in the stacks maximization (maxStacks) model [Boury et al, WABI 2026]

$$\text{Energy: } \#Stacks(R) = \sum_{\text{Helix } H \in R} |H| - 1$$

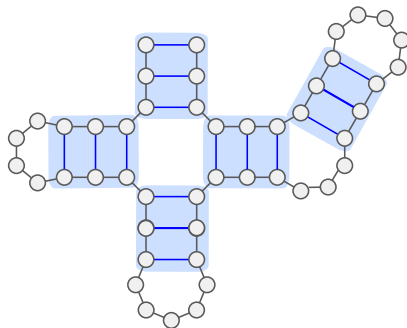
Definition (maxStacks inverse folding)

Input: Target secondary structure R

Output: RNA sequence $\omega \in \Sigma^*$ such that:

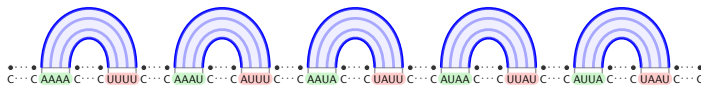
- ▶ Target R comp. with ω ;
- ▶ $\forall R' \neq R$, compatible with ω :

$$\#Stacks(R') < \#Stacks(R).$$



The maxStacks model differs from maxBP:

- ▶ Multiloops of arbitrary degree Δ can be designed



- ▶ **Result:** Exact linear algo (Locked⁺) for targets with helices greater than $\lceil \log_{3.56}(\Delta) \rceil + 6.2$ BPs

Stacks-maximizing designs are frequently designs for the Turner energy model

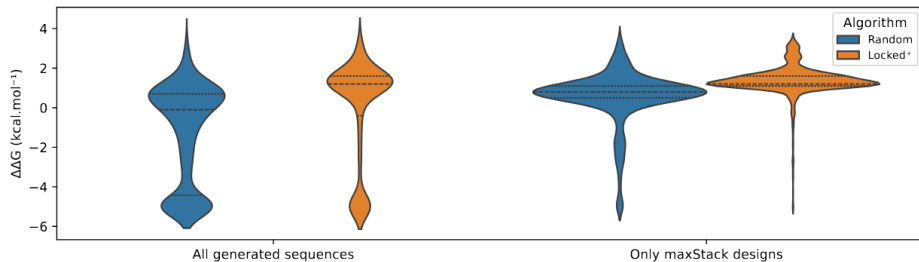
Definition: $\Delta\Delta G(\text{Candidate design } w, \text{Target } R) = \max_{R' \neq R} E(w; R') - E(w; R)$

$\Delta\Delta G \nearrow \rightarrow$ Greater stability relative to competitors

$\Delta\Delta G \leq 0 \rightarrow$ Dominated/challenged by competitor structure

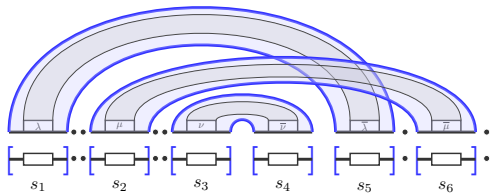
$\Delta\Delta G > 0 \rightarrow$ Solution to inverse folding

Random target structures: Length from 25 to 275nts, min helix length from 2 to 8 BPs



- ▶ Larger proportion of (better) Turner designs using Locked⁺ (maxStack algorithm)
- ▶ Stacking maximization usable as filter to improve potential as design

Inverse folding including crossing BPs/PseudoKnots follows for (almost) free



General **complexity open**, yet probably NP-Hard

Remark: General pseudoknots allowed **both** for input **and** alternative structures

Theorem (Boury et al, WABI 2026)

Linear-time exact algorithm for target structures with H (crossing) helices with more than $\lceil \log_{3.56}(H) \rceil + 6.2$ base pairs

Surprisingly, as checking if given RNA sequence represents solution is **NP-hard** [Lyngsø, ICALP'04]

Definition (maxStacks PK inv. folding)

Input: Target **crossing** structure R

Output: RNA sequence $\omega \in \Sigma^*$ such that:

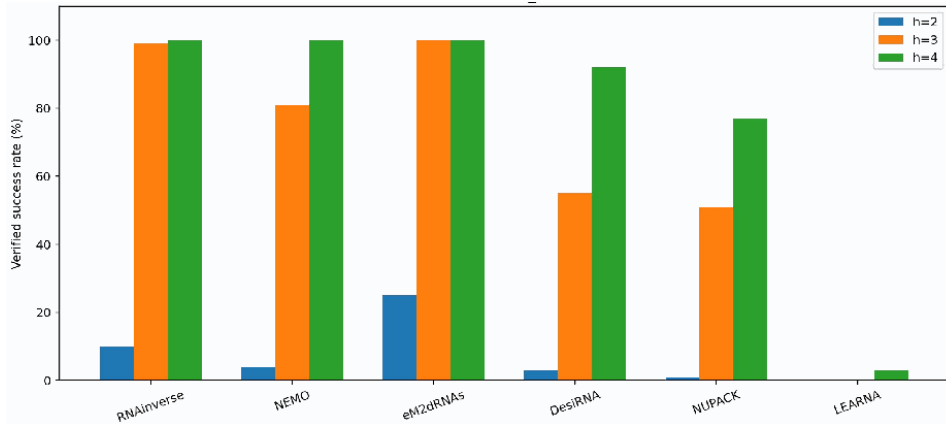
- ▶ Target R comp. with ω ;
- ▶ $\forall$ (crossing) $R' \neq R$, compatible with ω :
 $\#Stacks(R') < \#Stacks(R)$.

Conclusions

- ▶ Exact minimal solutions to inverse folding **improve empirical design performances**
- ▶ Longer helices **really** help structural design for deeper combinatorial reasons
- ▶ **Interaction design** with guaranteed specificity can be addressed in this framework
- ▶ Is RNA inverse folding **still hard**? Is method development **still relevant**?

Inverse folding remains hard anytime targets feature short helices

Random uniform target structure : Length 75 nts, min helix length 2 to 4

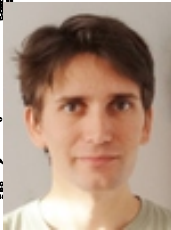


- ▶ Modest overall performances, especially in comparison to reported success rates
- ▶ Generalization issues of ML approaches?
- ▶ Old faithful – 3 decades old – RNAinverse still SOTA???

Thank you



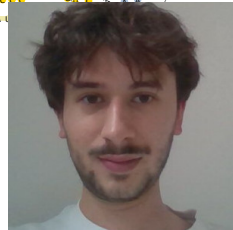
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- ▶ Ivo Hofacker and Leonhard Sidl (TBI Vienna)
- ▶ Sebastian Will and Sarah J. Berkemer (Ecole Polytechnique)
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